

Complexity Theory  
**Exercise 6: Diagonalization**  
8th December 2021

**Exercise 6.1.** Find the fault in the following proof of  $P \neq NP$ .

1. Assume that  $P = NP$ .
2. Then  $SAT \in P$  and thus there exists a  $k \in \mathbb{N}$  such that  $SAT \in DTIME(n^k)$ .
3. Because every language in  $NP$  is poly-time reducible to  $SAT$ , we have  $NP \subseteq DTIME(n^k)$ .
4. It follows that  $P \subseteq DTIME(n^k)$ .
5. By the Time Hierarchy Theorem there exist languages in  $DTIME(n^{k+1})$  that are not in  $DTIME(n^k)$ , contradicting  $P \subseteq DTIME(n^k)$ .
6. Therefore,  $P \neq NP$ .

**Exercise 6.2.** Show the following.

1.  $TIME(2^n) = TIME(2^{n+1})$
2.  $TIME_*(2^n) \subset TIME_*(2^{2n})$
3.  $NTIME(n) \subset PSPACE$

**Exercise 6.3.** Define a function that is computable but not time-constructible.

**Exercise 6.4.** Consider the function  $pad: \Sigma^* \times \mathbb{N} \rightarrow \Sigma^* \#^*$  defined as  $pad(s, \ell) = s \#^j$ , where  $j = \max(0, \ell - |s|)$ . For some language  $A \subseteq \Sigma^*$  and  $f: \mathbb{N} \rightarrow \mathbb{N}$  define  $pad(A, f) = \{ pad(s, f(|s|)) \mid s \in A \}$ .

Show all of the following statements.

1. Show that, if  $A \in DTIME(n^6)$ , then  $pad(A, n^2) \in DTIME(n^3)$ .
2. Show that, if  $NEXPTIME \neq EXPTIME$ , then  $P \neq NP$ .
3. Show for every  $A \subseteq \Sigma^*$  and  $k \in \mathbb{N}$  that  $A \in P$  if and only if  $pad(A, n^k) \in P$ .
4. Show  $P \neq DSPACE(n)$ .
5. Show  $NP \neq DSPACE(n)$ .

**Exercise 6.5.** You are given two oracles and one of them is the set **TRUE QBF**, but you do not know which one. Design a polynomial algorithm that decides **TRUE QBF** using these oracles.