Overview

1. Introduction | Relational data model
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3. Complexity of query answering
4. Complexity of FO query answering
5. Conjunctive queries
6. Tree-like conjunctive queries
7. Query optimisation
8. Conjunctive Query Optimisation / First-Order Expressiveness
9. First-Order Expressiveness / Introduction to Datalog
10. Expressive Power and Complexity of Datalog
11. Optimisation and Evaluation of Datalog
12. Evaluation of Datalog (2)
13. Graph Databases and Path Queries
14. Outlook: database theory in practice

See course homepage [⇒ link] for more information and materials
The relational queries considered so far produced a result table from a database. We generalize slightly.

**Definition**

- **Syntax:** a query expression $q$ is a word from a query language (algebra expression, logical expression, etc.)
- **Semantics:** a query mapping $M[q]$ is a function that maps a database instance $\mathcal{I}$ to a database instance $M[q](\mathcal{I})$

~~> a “result table” is a result database instance with one table.

~~> for some semantics, query mappings are not defined on all database instances
Generic Queries

We only consider queries that do not depend on the concrete names given to constants in the database:

**Definition**

A query $q$ is **generic** if, for every bijective renaming function $\mu : \text{dom} \to \text{dom}$ and database instance $\mathcal{I}$:

$$\mu(M[q](\mathcal{I})) = M[\mu(q)](\mu(\mathcal{I})).$$

In this case, $M[q]$ is **closed under isomorphisms**.
### Lines:

<table>
<thead>
<tr>
<th>Line</th>
<th>Type</th>
</tr>
</thead>
<tbody>
<tr>
<td>85</td>
<td>bus</td>
</tr>
<tr>
<td>3</td>
<td>tram</td>
</tr>
<tr>
<td>F1</td>
<td>ferry</td>
</tr>
<tr>
<td>...</td>
<td>...</td>
</tr>
</tbody>
</table>

### Stops:

<table>
<thead>
<tr>
<th>SID</th>
<th>Stop</th>
<th>Accessible</th>
</tr>
</thead>
<tbody>
<tr>
<td>17</td>
<td>Hauptbahnhof</td>
<td>true</td>
</tr>
<tr>
<td>42</td>
<td>Helmholtzstr.</td>
<td>true</td>
</tr>
<tr>
<td>57</td>
<td>Stadtgutstr.</td>
<td>true</td>
</tr>
<tr>
<td>123</td>
<td>Gustav-Freytag-Str.</td>
<td>false</td>
</tr>
<tr>
<td>...</td>
<td>...</td>
<td>...</td>
</tr>
</tbody>
</table>

### Connect:

<table>
<thead>
<tr>
<th>From</th>
<th>To</th>
<th>Line</th>
</tr>
</thead>
<tbody>
<tr>
<td>57</td>
<td>42</td>
<td>85</td>
</tr>
<tr>
<td>17</td>
<td>789</td>
<td>3</td>
</tr>
<tr>
<td>...</td>
<td>...</td>
<td>...</td>
</tr>
</tbody>
</table>

Every table has a **schema**:

- `Lines[Line:string, Type:string]`
- `Stops[SID:int, Stop:string, Accessible:bool]`
- `Connect[From:int, To:int, Line:string]`
First-order Logic as a Query Language

Idea: database instances are finite first-order interpretations
\rightsquigarrow use first-order formulae as query language
\rightsquigarrow use unnamed perspective (more natural here)

Examples (using schema as in previous lecture):

- Find all bus lines: \(\text{Lines}(x, "bus")\)
- Find all possible types of lines: \(\exists y.\text{Lines}(y, x)\)
- Find all lines that depart from an accessible stop:

\[
\exists y_{\text{SID}}, y_{\text{Stop}}, y_{\text{To}} . (\text{Stops}(y_{\text{SID}}, y_{\text{Stop}}, "true") \land \text{Connect}(y_{\text{SID}}, y_{\text{To}}, x_{\text{Line}}))
\]
First-order Logic with Equality: Syntax

Basic building blocks:

- **Predicate names** with an arity \( \geq 0 \): \( p, q \), Lines, Stops
- **Variables**: \( x, y, z \)
- **Constants**: \( a, b, c \)
- **Terms** are variables or constants: \( s, t \)

Formulae of first-order logic are defined as usual:

\[
\varphi ::= p(t_1, \ldots, t_n) \mid t_1 \approx t_2 \mid \neg \varphi \mid \varphi \land \varphi \mid \varphi \lor \varphi \mid \exists x. \varphi \mid \forall x. \varphi
\]

where \( p \) is an \( n \)-ary predicate, \( t_i \) are terms, and \( x \) is a variable.

- **An atom** is a formula of the form \( p(t_1, \ldots, t_n) \)
- **A literal** is an atom or a negated atom
- **Occurrences of variables** in the scope of a quantifier are **bound**; other occurrences of variables are **free**
First-order Logic Syntax: Simplifications

We use the usual shortcuts and simplifications:

- flat conjunctions ($\varphi_1 \land \varphi_2 \land \varphi_3$ instead of $(\varphi_1 \land (\varphi_2 \land \varphi_3))$)
- flat disjunctions (similar)
- flat quantifiers ($\exists x, y, z. \varphi$ instead of $\exists x. \exists y. \exists z. \varphi$)
- $\varphi \rightarrow \psi$ as shortcut for $\neg \varphi \lor \psi$
- $\varphi \leftrightarrow \psi$ as shortcut for $(\varphi \rightarrow \psi) \land (\psi \rightarrow \varphi)$
- $t_1 \not\approx t_2$ as shortcut for $\neg (t_1 \approx t_2)$

But we always use parentheses to clarify nesting of $\land$ and $\lor$:
No “$\varphi_1 \land \varphi_2 \lor \varphi_3$”!
First-order Logic with Equality: Semantics

First-order formulae are evaluated over interpretations $\langle \Delta^I, \cdot^I \rangle$, where $\Delta^I$ is the domain. To interpret formulas with free variables, we need a variable assignment $\mathcal{Z} : \text{Var} \rightarrow \Delta^I$.

- constants $a$ interpreted as $a^I, \mathcal{Z} = a^I \in \Delta^I$
- variables $x$ interpreted as $x^I, \mathcal{Z} = \mathcal{Z}(x) \in \Delta^I$
- $n$-ary predicates $p$ interpreted as $p^I \subseteq (\Delta^I)^n$
First-order Logic with Equality: Semantics

First-order formulae are evaluated over interpretations \( \langle \Delta^I,cdot^I \rangle \), where \( \Delta^I \) is the domain. To interpret formulas with free variables, we need a variable assignment \( Z : \text{Var} \rightarrow \Delta^I \).

- constants \( a \) interpreted as \( a^I,Z = a^I \in \Delta^I \)
- variables \( x \) interpreted as \( x^I,Z = Z(x) \in \Delta^I \)
- \( n \)-ary predicates \( p \) interpreted as \( p^I \subseteq (\Delta^I)^n \)

A formula \( \varphi \) can be satisfied by \( I \) and \( Z \), written \( I, Z \models \varphi \):

- \( I, Z \models p(t_1, \ldots, t_n) \) if \( \langle t_1^I,Z, \ldots, t_n^I,Z \rangle \in p^I \)
- \( I, Z \models t_1 \approx t_2 \) if \( t_1^I,Z = t_2^I,Z \)
- \( I, Z \models \neg \varphi \) if \( I, Z \nmodels \varphi \)
- \( I, Z \models \varphi \land \psi \) if \( I, Z \models \varphi \) and \( I, Z \models \psi \)
- \( I, Z \models \varphi \lor \psi \) if \( I, Z \models \varphi \) or \( I, Z \models \psi \)
- \( I, Z \models \exists x.\varphi \) if there is \( \delta \in \Delta^I \) with \( I, \{ x \mapsto \delta \}, Z \models \varphi \)
- \( I, Z \models \forall x.\varphi \) if for all \( \delta \in \Delta^I \) we have \( I, \{ x \mapsto \delta \}, Z \models \varphi \)
First-order Logic Queries

**Definition**

An *n*-ary first-order query $q$ is an expression $\varphi[x_1, \ldots, x_n]$ where $x_1, \ldots, x_n$ are exactly the free variables of $\varphi$ (in a specific order).

**Definition**

An **answer** to $q = \varphi[x_1, \ldots, x_n]$ over an interpretation $\mathcal{I}$ is a tuple $\langle a_1, \ldots, a_n \rangle$ of constants such that

$$\mathcal{I} \models \varphi[x_1/a_1, \ldots, x_n/a_n]$$

where $\varphi[x_1/a_1, \ldots, x_n/a_n]$ is $\varphi$ with each free $x_i$ replaced by $a_i$.

The **result** of $q$ over $\mathcal{I}$ is the set of all answers of $q$ over $\mathcal{I}$.
A Boolean query is a query of arity 0
⇒ we simply write $\varphi$ instead of $\varphi[]$
⇒ $\varphi$ is a closed formula (a.k.a. sentence)

What does a Boolean query return?
A Boolean query is a query of arity 0
\( \sim \) we simply write \( \varphi \) instead of \( \varphi[] \)
\( \sim \) \( \varphi \) is a closed formula (a.k.a. sentence)

What does a Boolean query return?

Two possible cases:

- \( \mathcal{I} \nmid \varphi \), then the result of \( \varphi \) over \( \mathcal{I} \) is \( \emptyset \) (the empty table)
- \( \mathcal{I} \models \varphi \), then the result of \( \varphi \) over \( \mathcal{I} \) is \( \{\langle\rangle\} \) (the unit table)

Interpreted as Boolean check with result true or false (match or no match)
Domain Dependence

We have defined FO queries over interpretations

How exactly do we get from databases to interpretations?

- Constants are just interpreted as themselves: $a^I = a$
- Predicates are interpreted according to the table contents
- But what is the domain of the interpretation?
Domain Dependence

We have defined FO queries over interpretations

How exactly do we get from databases to interpretations?

• Constants are just interpreted as themselves: \( a^I = a \)
• Predicates are interpreted according to the table contents
• But what is the domain of the interpretation?

What should the following queries return?

(1) \( \neg \text{Lines}(x, "bus")[x] \)
(2) \( (\text{Connect}(x_1, "42", "85") \lor \text{Connect}("57", x_2, "85"))[x_1, x_2] \)
(3) \( \forall y. p(x, y)[x] \)

Answers depend on the interpretation domain, not just on the database contents
First possible solution: the natural domain

Natural domain semantics (ND):

- fix the interpretation domain to $\text{dom}$ (infinite)
- query answers might be infinite (not a valid result table)
  $\leadsto$ query result undefined for such databases
Natural Domain: Examples

Query answers under natural domain semantics:

(1) \neg \text{Lines}(x, "bus")[x]
Natural Domain: Examples

Query answers under natural domain semantics:

(1) \( \neg \text{Lines}(x, "bus")[x] \)
    Undefined on all databases

(2) \((\text{Connect}(x_1, "42", "85") \lor \text{Connect}("57", x_2, "85"))[x_1, x_2]\)
Query answers under natural domain semantics:

(1) \( \neg \text{Lines}(x, \"bus\")[x] \)
   Undefined on all databases

(2) \( (\text{Connect}(x_1, \"42\", \"85\") \lor \text{Connect}(\"57\", x_2, \"85\"))[x_1, x_2] \)
   Undefined on databases with matching \( x_1 \) or \( x_2 \) in Connect, otherwise empty

(3) \( \forall y. p(x, y)[x] \)
Natural Domain: Examples

Query answers under natural domain semantics:

(1) \( \neg \text{Lines}(x, "bus")[x] \)
    Undefined on all databases

(2) \( (\text{Connect}(x_1, "42", "85") \lor \text{Connect}("57", x_2, "85"))[x_1, x_2] \)
    Undefined on databases with matching \( x_1 \) or \( x_2 \) in \text{Connect},
    otherwise empty

(3) \( \forall y. p(x, y)[x] \)
    Empty on all databases
Active Domain

Alternative: restrict to constants that are really used

→ active domain

- for a database instance \( I \), \( \text{adom}(I) \) is the set of constants used in relations of \( I \)
- for a query \( q \), \( \text{adom}(q) \) is the set of constants in \( q \)
- \( \text{adom}(I, q) = \text{adom}(I) \cup \text{adom}(q) \)

Active domain semantics (AD):
consider database instance as interpretation over \( \text{adom}(I, q) \)
Query answers under active domain semantics:

(1) ¬\text{Lines}(x, "bus")[x]
Query answers under active domain semantics:

(1) \( \neg \text{Lines}(x, "bus")[x] \)
    Let \( q' = \text{Lines}(x, "bus")[x] \). The answer is \( \text{adom}(I, q) \setminus M[q'](I) \)

(2) \( (\varphi_1[x_1] \lor \varphi_2[x_2]) \)[x_1, x_2]

Query answers under active domain semantics:

(1) \( \neg \text{Lines}(x, "bus") [x] \)

Let \( q' = \text{Lines}(x, "bus") [x] \). The answer is \( \text{adom}(I, q) \setminus M[q'](I) \)

(2) \( (\text{Connect}(x_1, "42", "85") \lor \text{Connect}("57", x_2, "85"))[x_1, x_2] \)

The answer is \( M[\varphi_1](I) \times \text{adom}(I, q) \cup \text{adom}(I, q) \times M[\varphi_2](I) \)

(3) \( \forall y. p(x, y)[x] \leadsto \text{see board} \)
Domain Independence

Observation: some queries do not depend on the domain

- \text{Stops}(x, y, "true")[x, y]
- \((x \approx a)[x]\)
- \(p(x) \land \neg q(x)[x]\)
- \(\forall y. (q(x, y) \rightarrow p(x, y))[x]\) (exercise: why?)

In contrast, all example queries on the previous few slides are not domain independent

Domain independent semantics (DI):
consider only domain independent queries
use any domain \(\text{adom}(I, q) \subseteq \Delta^I \subseteq \text{dom}\) for interpretation
How to Compare Query Languages

We have seen three ways of defining FO query semantics how to compare them?
How to Compare Query Languages

We have seen three ways of defining FO query semantics
\( \leadsto \) how to compare them?

**Definition**

The set of query mappings that can be described in a query language \( L \) is denoted \( \text{QM}(L) \).

- \( L_1 \) is **subsumed by** \( L_2 \), written \( L_1 \sqsubseteq L_2 \), if \( \text{QM}(L_1) \subseteq \text{QM}(L_2) \)
- \( L_1 \) is **equivalent to** \( L_2 \), written \( L_1 \equiv L_2 \), if \( \text{QM}(L_1) = \text{QM}(L_2) \)

We will also compare query languages under named perspective with query languages under unnamed perspective.
This is possible since there is an easy one-to-one correspondence between query mappings of either kind (see exercise).
Equivalence of Relational Query Languages

Theorem

The following query languages are equivalent:

- Relational algebra RA
- FO queries under active domain semantics AD
- Domain independent FO queries DI

This holds under named and under unnamed perspective.

To prove it, we will show:

\[
RA_{\text{named}} \subseteq DI_{\text{unnamed}} \subseteq AD_{\text{unnamed}} \subseteq RA_{\text{named}}
\]
$\text{RA}_{\text{named}} \equiv \text{DI}_{\text{unnamed}}$

For a given RA query $q[a_1, \ldots, a_n]$, we recursively construct a DI query $\varphi_q[x_{a_1}, \ldots, x_{a_n}]$ as follows:

We assume without loss of generality that all attribute lists in RA expressions respect the global order of attributes.

- if $q = R$ with signature $R[a_1, \ldots, a_n]$
RA\textsubscript{named} \sqsubseteq DI\textsubscript{unnamed}

For a given RA query \(q[a_1, \ldots, a_n]\), we recursively construct a DI query \(\varphi_q[x_{a_1}, \ldots, x_{a_n}]\) as follows:

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- if \(q = R\) with signature \(R[a_1, \ldots, a_n]\), then \(\varphi_q = R(x_{a_1}, \ldots, x_{a_n})\)

- if \(n = 1\) and \(q = \{a_1 \mapsto c\}\)
For a given RA query \( q[a_1, \ldots, a_n] \), we recursively construct a DI query \( \varphi_q[x_{a_1}, \ldots, x_{a_n}] \) as follows:

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- if \( q = R \) with signature \( R[a_1, \ldots, a_n] \), then \( \varphi_q = R(x_{a_1}, \ldots, x_{a_n}) \)
- if \( n = 1 \) and \( q = \{a_1 \mapsto c\} \), then \( \varphi_q = (x_{a_1} \approx c) \)
- if \( q = \sigma_{a_i=c}(q') \)
For a given RA query $q[a_1, \ldots, a_n]$, we recursively construct a DI query $\varphi_q[x_{a_1}, \ldots, x_{a_n}]$ as follows:

We assume without loss of generality that all attribute lists in RA expressions respect the global order of attributes.

- if $q = R$ with signature $R[a_1, \ldots, a_n]$, then $\varphi_q = R(x_{a_1}, \ldots, x_{a_n})$
- if $n = 1$ and $q = \{a_1 \mapsto c\}$, then $\varphi_q = (x_{a_1} \approx c)$
- if $q = \sigma_{a_i=c}(q')$, then $\varphi_q = \varphi_{q'} \land (x_{a_i} \approx c)$
- if $q = \sigma_{a_i=a_j}(q')$
For a given RA query \( q[a_1, \ldots, a_n] \), we recursively construct a DI query \( \varphi_q[x_{a_1}, \ldots, x_{a_n}] \) as follows:

We assume without loss of generality that all attribute lists in RA expressions respect the global order of attributes.

- if \( q = R \) with signature \( R[a_1, \ldots, a_n] \), then \( \varphi_q = R(x_{a_1}, \ldots, x_{a_n}) \)

- if \( n = 1 \) and \( q = \{\{a_1 \mapsto c\}\} \), then \( \varphi_q = (x_{a_1} \approx c) \)

- if \( q = \sigma_{a_i=c}(q') \), then \( \varphi_q = \varphi_{q'} \land (x_{a_i} \approx c) \)

- if \( q = \sigma_{a_i=a_j}(q') \), then \( \varphi_q = \varphi_{q'} \land (x_{a_i} \approx x_{a_j}) \)

- if \( q = \delta_{b_1,\ldots,b_n\rightarrow a_1,\ldots,a_n} q' \)
RA_{\text{named}} \sqsubseteq \text{DI}_{\text{unnamed}}

For a given RA query $q[a_1, \ldots, a_n]$, we recursively construct a DI query $\varphi_q[x_{a_1}, \ldots, x_{a_n}]$ as follows:

We assume without loss of generality that all attribute lists in RA expressions respect the global order of attributes.

- if $q = R$ with signature $R[a_1, \ldots, a_n]$, then $\varphi_q = R(x_{a_1}, \ldots, x_{a_n})$
- if $n = 1$ and $q = \{a_1 \mapsto c\}$, then $\varphi_q = (x_{a_1} \approx c)$
- if $q = \sigma_{a_i=c}(q')$, then $\varphi_q = \varphi_{q'} \land (x_{a_i} \approx c)$
- if $q = \sigma_{a_i=a_j}(q')$, then $\varphi_q = \varphi_{q'} \land (x_{a_i} \approx x_{a_j})$
- if $q = \delta_{b_1, \ldots, b_n \rightarrow a_1, \ldots, a_n} q'$, then

$$\varphi_q = \exists y_{b_1}, \ldots, y_{b_n}. (x_{a_1} \approx y_{b_1}) \land \ldots \land (x_{a_n} \approx y_{b_n}) \land \varphi'_{q'[y_{a_1}, \ldots, y_{a_n}]}$$

(Here we assume that the $a_1, \ldots, a_n$ in $\delta_{b_1, \ldots, b_n \rightarrow a_1, \ldots, a_n}$ are written in the order of attributes, whereas $b_1, \ldots, b_n$ might be in another order. $\varphi'_{q'[y_{a_1}, \ldots, y_{a_n}]}$ is like $\varphi_{q'}$ but using variables $y_{a_i}$.)
$$\text{RA}_{\text{named}} \equiv \text{DI}_{\text{unnamed}} \ (\text{cont'd})$$

Remaining cases:

- if $q = \pi_{a_1, \ldots, a_n}(q')$ for a subquery $q'[b_1, \ldots, b_m]$ with

$$\{b_1, \ldots, b_m\} = \{a_1, \ldots, a_n\} \cup \{c_1, \ldots, c_k\}$$
Remaining cases:

- if $q = \pi_{a_1, \ldots, a_n}(q')$ for a subquery $q'[b_1, \ldots, b_m]$ with $\{b_1, \ldots, b_m\} = \{a_1, \ldots, a_n\} \cup \{c_1, \ldots, c_k\}$, then $\varphi_q = \exists x_{c_1}, \ldots, x_{c_k} \cdot \varphi_{q'}$

- if $q = q_1 \bowtie q_2$
RA_{\text{named}} \sqsubseteq DI_{\text{Unnamed}} \ (\text{cont'd})

Remaining cases:

- if \( q = \pi_{a_1, \ldots, a_n}(q') \) for a subquery \( q'[b_1, \ldots, b_m] \) with
  \[ \{b_1, \ldots, b_m\} = \{a_1, \ldots, a_n\} \cup \{c_1, \ldots, c_k\}, \]
  then \( \varphi_q = \exists x_{c_1}, \ldots, x_{c_k}. \varphi_{q'} \)

- if \( q = q_1 \bowtie q_2 \) then \( \varphi_q = \varphi_{q_1} \land \varphi_{q_2} \)

- if \( q = q_1 \cup q_2 \)
RA_{named} \equiv DI_{unnamed} \ (cont'd)

Remaining cases:

- if $q = \pi_{a_1,\ldots,a_n}(q')$ for a subquery $q'[b_1,\ldots,b_m]$ with
  
  \{b_1,\ldots,b_m\} = \{a_1,\ldots,a_n\} \cup \{c_1,\ldots,c_k\},

  then $\varphi_q = \exists x_{c_1},\ldots,x_{c_k} \cdot \varphi_{q'}$

- if $q = q_1 \bowtie q_2$ then $\varphi_q = \varphi_{q_1} \land \varphi_{q_2}$

- if $q = q_1 \cup q_2$ then $\varphi_q = \varphi_{q_1} \lor \varphi_{q_2}$

- if $q = q_1 - q_2$
RA\textsubscript{named} \sqsubseteq DI\textsubscript{unnamed} (cont’d)

Remaining cases:

- if $q = \pi_{a_1, \ldots, a_n}(q')$ for a subquery $q'[b_1, \ldots, b_m]$ with
  $$\{b_1, \ldots, b_m\} = \{a_1, \ldots, a_n\} \cup \{c_1, \ldots, c_k\},$$
  then $\varphi_q = \exists x_{c_1}, \ldots, x_{c_k} \cdot \varphi_q'$

- if $q = q_1 \bowtie q_2$ then $\varphi_q = \varphi_{q_1} \land \varphi_{q_2}$

- if $q = q_1 \cup q_2$ then $\varphi_q = \varphi_{q_1} \lor \varphi_{q_2}$

- if $q = q_1 - q_2$ then $\varphi_q = \varphi_{q_1} \land \neg \varphi_{q_2}$

One can show that $\varphi_q[x_{a_1}, \ldots, x_{a_n}]$ is domain independent and equivalent to $q$

$\sim$ exercise
$\text{DI}_{\text{unnamed}} \sqsubseteq \text{AD}_{\text{unnamed}}$

This is easy to see
This is easy to see:

- Consider an FO query $q$ that is domain independent
- The semantics of $q$ is the same for any domain $\text{adom} \subseteq \Delta^T \subseteq \text{dom}$
- In particular, the semantics of $q$ is the same under active domain semantics
- Hence, for every DI query, there is an equivalent AD query
Consider an AD query \( q = \varphi[x_1, \ldots, x_n] \).

For an arbitrary attribute name \( a \), we can construct an RA expression \( E_{a, \text{adom}} \) such that \( E_{a, \text{adom}}(I) = \{a \mapsto c \mid c \in \text{adom}(I, q)\} \)

\( \sim \) exercise
\[ \text{AD}_{\text{unnamed}} \sqsubseteq \text{RA}_{\text{named}} \]

Consider an AD query \( q = \varphi[x_1, \ldots, x_n] \).

For an arbitrary attribute name \( a \), we can construct an RA expression \( E_{a, \text{adom}} \) such that \( E_{a, \text{adom}}(\mathcal{I}) = \{a \mapsto c \mid c \in \text{adom}(\mathcal{I}, q)\} \).

Exercise

For every variable \( x \), we use a distinct attribute name \( a_x \)

- if \( \varphi = R(t_1, \ldots, t_m) \) with signature \( R[a_1, \ldots, a_m] \) with variables \( x_1 = t_{v_1}, \ldots, x_n = t_{v_n} \) and constants \( c_1 = t_{w_1}, \ldots, c_k = t_{w_k} \).
Consider an AD query $q = \varphi[x_1, \ldots, x_n]$. 

For an arbitrary attribute name $a$, we can construct an RA expression $E_{a, \text{adom}}$ such that $E_{a, \text{adom}}(\mathcal{I}) = \{\{a \mapsto c\} | c \in \text{adom}(\mathcal{I}, q)\}$. 

**exercise**

For every variable $x$, we use a distinct attribute name $a_x$.

- if $\varphi = R(t_1, \ldots, t_m)$ with signature $R[a_1, \ldots, a_m]$ with variables $x_1 = t_{v_1}, \ldots, x_n = t_{v_n}$ and constants $c_1 = t_{w_1}, \ldots, c_k = t_{w_k}$, then $E_\varphi = \delta_{a_{v_1} \ldots a_{v_n}} \rightarrow a_{x_1} \ldots a_{x_n} (\sigma_{a_{w_1} = c_1} (\ldots \sigma_{a_{w_k} = c_k} (R) \ldots))$

- if $\varphi = (x \approx c)$
Consider an AD query $q = \varphi[x_1, \ldots, x_n]$.

For an arbitrary attribute name $a$, we can construct an RA expression $E_{a,\text{dom}}$ such that $E_{a,\text{dom}}(I) = \{a \mapsto c \mid c \in \text{dom}(I, q)\}$

For every variable $x$, we use a distinct attribute name $a_x$

- if $\varphi = R(t_1, \ldots, t_m)$ with signature $R[a_1, \ldots, a_m]$ with variables $x_1 = t_{v_1}, \ldots, x_n = t_{v_n}$ and constants $c_1 = t_{w_1}, \ldots, c_k = t_{w_k}$, then $E_{\varphi} = \delta_{a_{v_1} \ldots a_{v_n} \rightarrow a_1 \ldots a_n} (\sigma_{a_{w_1} = c_1} (\ldots \sigma_{a_{w_k} = c_k} (R) \ldots))$

- if $\varphi = (x \approx c)$, then $E_{\varphi} = \{a_x \mapsto c\}$

- if $\varphi = (x \approx y)$
Consider an AD query \( q = \varphi[x_1, \ldots, x_n] \).

For an arbitrary attribute name \( a \), we can construct an RA expression \( E_{a, \text{adom}} \) such that \( E_{a, \text{adom}}(I) = \{\{a \mapsto c\} \mid c \in \text{adom}(I, q)\} \)

For every variable \( x \), we use a distinct attribute name \( a_x \)

- if \( \varphi = R(t_1, \ldots, t_m) \) with signature \( R[a_1, \ldots, a_m] \) with variables \( x_1 = t_{v_1}, \ldots, x_n = t_{v_n} \) and constants \( c_1 = t_{w_1}, \ldots, c_k = t_{w_k} \), then \( E\varphi = \delta_{a_{v_1} \ldots a_{v_n} \rightarrow a_{x_1} \ldots a_{x_n}} (\sigma_{a_{w_1} = c_1} (\ldots \sigma_{a_{w_k} = c_k} (R) \ldots)) \)

- if \( \varphi = (x \approx c) \), then \( E\varphi = \{a_x \mapsto c\} \)

- if \( \varphi = (x \approx y) \), then \( E\varphi = \sigma_{a_x = a_y}(E_{a_x, \text{adom}} \Join E_{a_y, \text{adom}}) \)

- other forms of equality atoms are similar
AD_{unamed} \equiv RA_{named} (cont’d)

Remaining cases:

- if $\varphi = \neg \psi$

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AD_{unnamed} \subseteq RA_{named} (cont’d)

Remaining cases:

- if $\varphi = \neg \psi$, then $E\varphi = (E_{a_{x_1}, \text{dom}} \bowtie \ldots \bowtie E_{a_{x_n}, \text{dom}}) - E\psi$
- if $\varphi = \varphi_1 \land \varphi_2$
Remaining cases:

- if $\varphi = \neg \psi$, then $E\varphi = (E_{a_1, \text{dom}} \bowtie \ldots \bowtie E_{a_n, \text{dom}}) - E\psi$

- if $\varphi = \varphi_1 \land \varphi_2$, then $E\varphi = E\varphi_1 \bowtie E\varphi_2$

- if $\varphi = \exists y. \psi$ where $\psi$ has free variables $y, x_1, \ldots, x_n$
AD_{unnamed} \equiv RA_{named} \ (cont’d)

Remaining cases:

- if $\varphi = \neg \psi$, then $E\varphi = (E_{a_{x_1}, \text{dom}} \Join \ldots \Join E_{a_{x_n}, \text{dom}}) - E\psi$

- if $\varphi = \varphi_1 \land \varphi_2$, then $E\varphi = E\varphi_1 \Join E\varphi_2$

- if $\varphi = \exists y. \psi$ where $\psi$ has free variables $y, x_1, \ldots, x_n$, then $E\varphi = \pi_{a_{x_1}, \ldots, a_{x_n}} E\psi$

The cases for $\lor$ and $\forall$ can be constructed from the above.

$\rightarrow$ exercise
\[ \text{AD}_{\text{unnamed}} \sqsubseteq \text{RA}_{\text{named}} \text{ (cont'd)} \]

Remaining cases:

- if \( \varphi = \neg \psi \), then \( E\varphi = (E_{a_1, \text{dom}} \bowtie \ldots \bowtie E_{a_n, \text{dom}}) - E\psi \)

- if \( \varphi = \varphi_1 \land \varphi_2 \), then \( E\varphi = E\varphi_1 \bowtie E\varphi_2 \)

- if \( \varphi = \exists y. \psi \) where \( \psi \) has free variables \( y, x_1, \ldots, x_n \), then \( E\varphi = \pi_{a_{x_1}, \ldots, a_{x_n}} E\psi \)

The cases for \( \lor \) and \( \forall \) can be constructed from the above exercise

A note on order: The translation yields an expression \( E\varphi[a_{x_1}, \ldots, a_{x_n}] \). For this to be equivalent to the query \( \varphi[x_1, \ldots, x_n] \), we must choose the attribute names such that their global order is \( a_{x_1}, \ldots, a_{x_n} \). This is clearly possible, since the names are arbitrary and we have infinitely many names available.
How to find DI queries?

Domain independent queries are arguably most intuitive, since their result does not depend on special assumptions.

How can we check if a query is in DI?
How to find DI queries?

Domain independent queries are arguably most intuitive, since their result does not depend on special assumptions.

\( \Rightarrow \) How can we check if a query is in DI? Unfortunately, we can’t:

**Theorem**

Given a FO query \( q \), it is undecidable if \( q \in \text{DI} \).

\( \Rightarrow \) find decidable sufficient conditions for a query to be in DI
We first define a normal form for FO queries:

**Safe-Range Normal Form (SRNF)**

- Rename variables apart (distinct quantifiers bind distinct variables, bound variables distinct from free variables)
- Eliminate all universal quantifiers: $\forall y.\psi \mapsto \neg \exists y.\neg \psi$
- Push negations inwards:
  - $\neg(\varphi \land \psi) \mapsto (\neg \varphi \lor \neg \psi)$
  - $\neg(\varphi \lor \psi) \mapsto (\neg \varphi \land \neg \psi)$
  - $\neg\neg \psi \mapsto \psi$
Safe-Range Queries

Let $\varphi$ be a formula in SRNF. The set $\text{rr}(\varphi)$ of range-restricted variables of $\varphi$ is defined recursively:

$$\text{rr}(R(t_1, \ldots, t_n)) = \{ x \mid x \text{ a variable among the } t_1, \ldots, t_n \}$$

$$\text{rr}(x \approx a) = \{ x \}$$

$$\text{rr}(x \approx y) = \emptyset$$

$$\text{rr}(\varphi_1 \land \varphi_2) = \begin{cases} 
\text{rr}(\varphi_1) \cup \{ x, y \} & \text{if } \varphi_2 = (x \approx y) \text{ and } \{ x, y \} \cap \text{rr}(\varphi_1) \neq \emptyset \\
\text{rr}(\varphi_1) \cup \text{rr}(\varphi_2) & \text{otherwise}
\end{cases}$$

$$\text{rr}(\varphi_1 \lor \varphi_2) = \text{rr}(\varphi_1) \cap \text{rr}(\varphi_2)$$

$$\text{rr}(\exists y. \psi) = \begin{cases} 
\text{rr}(\psi) \setminus \{ y \} & \text{if } y \in \text{rr}(\psi) \\
\text{throw new NotSafeException()} & \text{if } y \notin \text{rr}(\psi)
\end{cases}$$

$$\text{rr}(\neg \psi) = \emptyset \text{ if } \text{rr}(\psi) \text{ is defined (no exception)}$$
Safe-Range Queries

**Definition**

An FO query $q = \varphi[x_1, \ldots, x_n]$ is a **safe-range query** if

$$\text{rr}(\text{SRNF}(\varphi)) = \{x_1, \ldots, x_n\}.$$  

Safe-range queries are domain independent.
**Definition**

An FO query $q = \varphi[x_1, \ldots, x_n]$ is a **safe-range query** if

$$\text{rr}(\text{SRNF}(\varphi)) = \{x_1, \ldots, x_n\}.$$ 

Safe-range queries are domain independent.
One can show a much stronger result:

**Theorem**

The following query languages are equivalent:

- Safe-range queries SR
- Relational algebra RA
- FO queries under active domain semantics AD
- Domain independent FO queries DI
There are more equivalent ways to define a relational query language.

Example: Codd’s tuple calculus

- Based on named perspective
- Use first-order logic, but variables range over sorted tuples (rows) instead of values
- Use expressions like $x: \text{From,To,Line}$ to declare sorts of variables in queries
- Use expressions like $x.\text{From}$ to access a specific value of a tuple
- Example: Find all lines that depart from an accessible stop

\[
\{x: \text{Line} \mid \exists y: \text{SID,Stop,Accessible}. (\text{Stops}(y) \land y.\text{Accessible} \approx \text{"true"} \land \exists z: \text{From,To,Line}. (\text{Connect}(z) \land z.\text{From} \approx y.\text{SID} \land z.\text{Line} \approx x.\text{Line}) \})\}
\]
First-order logic gives rise to a relational query language

The problem of domain dependence can be solved in several ways

All common definitions lead to equivalent calculi

$\leadsto$ “relational calculus”

Open questions:

- How hard is it to actually answer such queries? (next lecture)
- How can we study the expressiveness of query languages?
- Are there interesting query languages that are not equivalent to RA?