TECHNISCHE
UNIVERSITÄT
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DATABASE THEORY

Lecture 11: Query Expressiveness

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TU Dresden, 13 May 2025

Review

First-Order Query Expressiveness

Queries and Their Expressiveness

Recall:

- Syntax: a query expression q is a word from a query language (algebra expression, logical expression, etc.)
- Semantics: a query mapping M is a function that maps a database instance I to a database table $M(I)$
- We only study generic queries, which are closed under bijective renaming (isomorphism of databases)

> **Definition 11.1:** The expressiveness of a query language is characterised by the set of query mappings that it can express.
>
> Given a query language **L**, a query mapping M is **L**-definable if there is a query expression $q \in$ **L** such that $M[q] = M$.

We can study expressiveness for all query mappings over all possible databases, or we can restrict attention to a subset of query mappings or to a subset of databases.

Boolean Query Mappings

A Boolean query mapping is a query mapping that returns "true" (usually a table with one empty row) or "false" (usually an empty table).
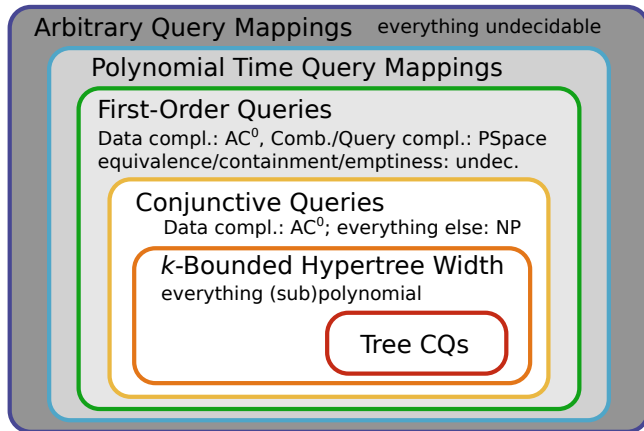
Every Boolean query mapping

- defines set of databases for which it is true
- defines a decision problem over the set of all databases
- could be decidable or undecidable
- if decidable, it may be characterised in terms of complexity

  Note: the "complexity of a mapping" is always "data complexity," i.e., complexity w.r.t. the size of the input database; the mapping defines the decision problem and is fixed.

Expressivity vs. Complexity

All query mappings that can be expressed in first-order logic are of polynomial complexity, actually in AC^0.

Arbitrary Query Mappings everything undecidable

Polynomial Time Query Mappings

First-Order Queries
Data compl.: AC^0, Comb./Query compl.: PSpace
equivalence/containment/emptiness: undec.

Conjunctive Queries
Data compl.: AC^0; everything else: NP

k-Bounded Hypertree Width
everything (sub)polynomial

Tree CQs

The Limits of FO Queries

Are there polynomial query mappings that cannot be expressed in FO?
$\leadsto$ yes!

We already knew this from previous lectures:

- We learned that $AC^0 \subset NC^1 \subseteq \ldots \subseteq P$
- Hence, there is a problem X in NC^1 that is not in AC^0
- Therefore, the corresponding query mapping M_X is not FO-definable

$AC^0 \subset NC^1$ was first shown for the problem $X = \text{PARITY}$:

- **Input:** finite relational structure $\mathcal{I}$
- **Output:** "true" if $\mathcal{I}$ has an even number of domain elements

The original proof is specific to this problem [Ajtai 1983].

Any Other FO-Undefinable Problems?

Yes, many.

Strong evidence from complexity theory:

- If any P-complete problem X were FO-definable,
- then every problem in P could be LogSpace-reduced to X
- and then solved in AC^0,
- hence every problem in P could be solved in LogSpace,
- that is, P = L.
- Most experts do not think that this is the case.

Therefore, one would expect all P-hard and similarly all NL-hard problems to not be FO-definable.

$\rightsquigarrow$ How can we see this more directly?

Proving FO-Undefinability

How to show that a query mapping is FO-definable?
↝ Find an FO query that expresses the query mapping

How to show that a query mapping is **not** FO-definable?
↝ Not so easy ... important tools:

- Ehrenfeucht-Fraïssé games
- Locality theorems

Ehrenfeucht-Fraïssé Games

A method for showing that certain finite structures cannot be distinguished by certain FO formulas

General idea:

- A game is played on two databases $\mathcal{I}$ and $\mathcal{J}$
- There are two players: the Spoiler and the Duplicator
- The players select elements from $\mathcal{I}$ and $\mathcal{J}$ in each round
- Spoiler wants to show that the two databases are different
- Duplicator wants make the databases appear to be the same

We will always play on finite structures without constant symbols
(remember that one can simulate constants by unary relations with one row)

Playing One Run of an EF Game

A single run of the game has a fixed number r of rounds

Spoiler starts each round, and Duplicator answers:

- Spoiler picks a domain element from $\mathcal{I}$ or from $\mathcal{J}$
- Duplicator picks an element from the other database ($\mathcal{J}$ or $\mathcal{I}$)

$\rightsquigarrow$ One element gets picked from each $\mathcal{I}$ and $\mathcal{J}$ per round

$\rightsquigarrow$ Run of game ends with two lists of elements:
$a_1, \ldots, a_r \in \Delta^{\mathcal{I}}$ and $b_1, \ldots, b_r \in \Delta^{\mathcal{J}}$

Duplicator wins the run if:

- For all indices i and j, we have $a_i = a_j$ if and only if $b_i = b_j$.
- For all lists of indices $i_1, \ldots, i_n$ and n-ary relation names R,
  we have $\langle a_{i_1}, \ldots, a_{i_n} \rangle \in R^{\mathcal{I}}$ if and only if $\langle b_{i_1}, \ldots, b_{i_n} \rangle \in R^{\mathcal{J}}$.

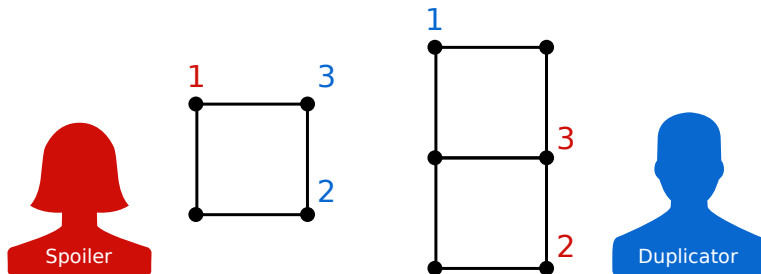 "The substructures induced by the selected elements are isomorphic"

Otherwise Spoiler wins the run.

- edges denote a bi-directional binary predicate
- all edges are the same predicate
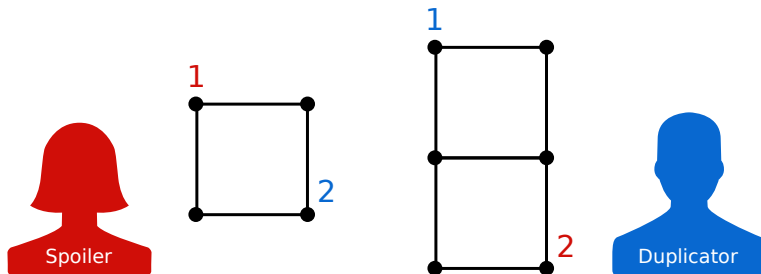
- edges denote a bi-directional binary predicate
- all edges are the same predicate

Winning the EF Game

The game is won by whoever has a winning strategy:

> A player has a winning strategy if they can make sure that they will win, whatever the other player is doing.

In other words:

- Duplicator wins if he can duplicate any move that the spoiler makes.
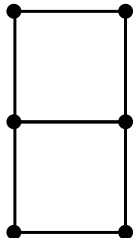- Spoiler wins if she can spoil any attempt to duplicate her moves.
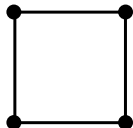
We write $\mathcal{I} \sim_r \mathcal{J}$ if Duplicator wins the r-round EF game on $\mathcal{I}$ and $\mathcal{J}$.

Observation: given enough moves, the spoiler will always win, unless the structures are isomorphic

Example

Who wins the 2-round game?
Who wins the 3-round game?

- edges denote a bi-directional binary predicate
- all edges are the same predicate

Quantifier Rank

EF games characterise expressivity of FO formulae based on the nesting depth of quantifiers:

Definition 11.2: The quantifier rank of a FO formula is the maximal nesting level of quantifiers within the formula.

Example 11.3:

- A formula without quantifiers has quantifier rank 0
- $\exists x.(C(x) \land \forall y.(R(x, y) \rightarrow x \approx y) \land \exists v.S(x, v))$ has quantifier rank 2

Definition 11.4: We write $\mathcal{I} \equiv_r \mathcal{J}$ if $\mathcal{I}$ and $\mathcal{J}$ satisfy the same FO sentences of rank r (or less).

Significance of EF Games

Theorem 11.5: For every r, $\mathcal{I}$ and $\mathcal{J}$, the following are equivalent:

- $\mathcal{I} \equiv_r \mathcal{J}$, that is, $\mathcal{I}$ and $\mathcal{J}$ satisfy the same FO sentences of rank r (or less).
- $\mathcal{I} \sim_r \mathcal{J}$, that is, the Duplicator wins the r-round EF game on $\mathcal{I}$ and $\mathcal{J}$.

Therefore, the following are equivalent:

- The query mapping M is FO-definable
- There is an FO sentence φ that defines M
- There is a number r such that, for every $\mathcal{I}$ accepted by M and every $\mathcal{J}$ not accepted by M, the Spoiler wins the r-round EF game on $\mathcal{I}$ and $\mathcal{J}$

Proof idea (1)

We outline the proof for the direction that is more important to us:

> **Lemma 11.6:** For every r, we find $\sim_r \subseteq \equiv_r$.

Proof: We show the contrapositive: if $\mathcal{I} \not\equiv_r \mathcal{J}$ then $\mathcal{I} \not\sim_r \mathcal{J}$. Hence, suppose there is a formula φ_r of quantifier depth r such that (w.l.o.g.) $\mathcal{I} \models \varphi_r$ and $\mathcal{J} \models \neg\varphi_r$.

We sketch the idea for the case that φ_r is in prenex normal form $\varphi_r = Q_1 x_1 \ldots Q_r.x_r.\psi$ with $Q_i \in \{\exists, \forall\}$ and ψ a quantifier-free formula:

- Then $\neg\varphi_r$ is equivalent to $\bar{Q}_1 x_1 \ldots \bar{Q}_r.x_r.\neg\psi$, where $\bar{\exists} = \forall$ and $\bar{\forall} = \exists$
- Spoiler will enforce a selection of elements $a_1, \ldots, a_r \in \Delta^{\mathcal{I}}$ and $b_1, \ldots, b_r \in \Delta^{\mathcal{J}}$, such that, after i steps of the game, $\mathcal{I}, \{x_1 \mapsto a_1, \ldots, x_i \mapsto a_i\} \models Q_{i+1} x_{i+1} \ldots Q_r x_r.\psi$ and $\mathcal{J}, \{x_1 \mapsto b_1, \ldots, x_i \mapsto b_i\} \not\models Q_{i+1} x_{i+1} \cdots Q_r x_r.\psi$ ($*$):
    - Property ($*$) holds initially ($i = 0$) by assumption.
    - In step $i + 1$, if $Q_{i+1} = \exists$, Spoiler selects $a_{i+1} \in \Delta^{\mathcal{I}}$ such that $\mathcal{I}, \{x_1 \mapsto a_1, \ldots, x_{i+1} \mapsto a_{i+1}\} \models Q_{i+2} x_{i+2} \ldots Q_r x_r.\psi$ – this exists because of ($*$).
    - Any choice b_{i+1} of Duplicator will be such that $\mathcal{J}, \{x_1 \mapsto b_1, \ldots, x_{i+1} \mapsto b_{i+1}\} \not\models Q_{i+2} x_{i+2} \ldots Q_r x_r.\psi$, since $\bar{Q}_{i+1} = \forall$.
    - The case $Q_{i+1} = \forall$ is similar: now Spoiler selects b_{i+1}.

Proof idea (2)

We outline the proof for the direction that is more important to us:

Lemma 11.6: For every r, we find $\sim_r \,\subseteq\, \equiv_r$.

Proof (continued): Therefore, by $(*)$, after r rounds we have selected elements $a_1, \ldots, a_r \in \Delta^{\mathcal{I}}$ and $b_1, \ldots, b_r \in \Delta^{\mathcal{J}}$, such that $\mathcal{I}, \{x_1 \mapsto a_1, \ldots, x_r \mapsto a_r\} \models \psi$ and $\mathcal{J}, \{x_1 \mapsto b_1, \ldots, x_r \mapsto b_r\} \not\models \psi$.

Hence, the substructures induced by the selected elements are not isomorphic (if they were, we would find that ψ evaluates to the same in both cases)

$\rightsquigarrow$ Spoiler wins

The idea can be generalised to formulae φ_r that are not in prenex normal form (by interleaving the choice of the quantifier and the evaluation of the formula) $\qquad\square$

Example

Let's assume all edges denote the (bi-directional) predicate r:
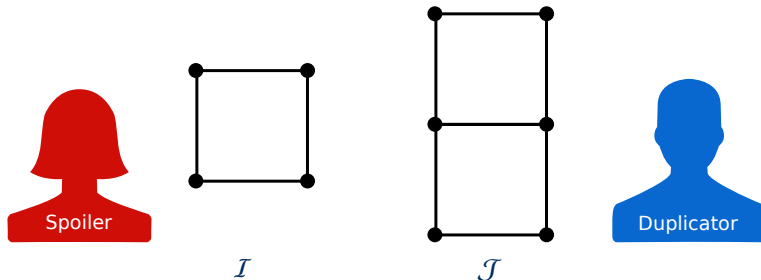
Which formula distinguishes the two structures?

For example: $\varphi_3 = \exists x.\exists y.\forall z.r(x,z) \leftrightarrow r(y,z)$

- $\mathcal{I} \models \varphi_3$
- $\mathcal{J} \not\models \varphi_3$

The formula corresponds to a 3-move winning strategy for Spoiler:

- first select opposing corners in $\mathcal{I}$
- then select an element in $\mathcal{J}$ that neighbours exactly one of the elements selected by Duplicator

Using EF Games to Show FO-Undefinability

How to show that a query mapping M can not be FO-defined:

- Let C_M be the class of all databases recognised by M
- Find sequences of databases $\mathcal{I}_1, \mathcal{I}_2, \mathcal{I}_3, \ldots \in C_M$ and databases $\mathcal{J}_1, \mathcal{J}_2, \mathcal{J}_3, \ldots \notin C_M$, such that $\mathcal{I}_i \sim_i \mathcal{J}_i$

$\rightsquigarrow$ for any formula φ (however large its quantifier rank r), there is a counterexample $\mathcal{I}_r \in C_M$ and $\mathcal{J}_r \notin C_M$ that φ cannot distinguish

Problems:

- How to find such sequences of $\mathcal{I}_i$ and $\mathcal{J}_i$?
  - $\rightsquigarrow$ No general strategy exists
- Given suitable sequences, how to show that $\mathcal{I}_i \sim_i \mathcal{J}_i$?
  - $\rightsquigarrow$ Can be difficult, but doable for some special cases

Expressiveness on Linear Orders

Let's look at some very simple structures:

> **Definition 11.7:** A structure $\mathcal{I}$ is a linear order if it has a single binary predicate $\leq$ interpreted as a total, transitive, reflexive and asymmetric relation.

> **Example 11.8:** Consider the following structures:
>
> $\mathcal{L}_6 : 1 \leq 2 \leq 3 \leq 4 \leq 5 \leq 6$
> $\mathcal{L}_7 : 1 \leq 2 \leq 3 \leq 4 \leq 5 \leq 6 \leq 7$
>
> Who wins the 3-round EF game? Spoiler can win the 3-round EF game as follows:
>
> Spoiler plays 4 in $\mathcal{L}_7$
>   Duplicator plays 4 in $\mathcal{L}_6$: Spoiler plays 6 in $\mathcal{L}_7$
>   Duplicator plays 5 in $\mathcal{L}_6$: Spoiler plays 5 in $\mathcal{L}_7$ and wins
>   Duplicator plays 6 in $\mathcal{L}_6$: Spoiler plays 7 in $\mathcal{L}_7$ and wins
>   Duplicator plays 3 in $\mathcal{L}_6$: symmetric game (flipped horizontally)

Expressiveness on Linear Orders

Let's look at some very simple structures:

Definition 11.7: A structure $\mathcal{I}$ is a linear order if it has a single binary predicate $\leq$ interpreted as a total, transitive, reflexive and asymmetric relation.

Example 11.9: Consider the following structures:

$\mathcal{L}_7 : 1 \leq 2 \leq 3 \leq 4 \leq 5 \leq 6 \leq 7$
$\mathcal{L}_8 : 1 \leq 2 \leq 3 \leq 4 \leq 5 \leq 6 \leq 7 \leq 8$

Spoiler cannot win the 3-round EF game:

Spoiler plays 4 in $\mathcal{L}_8$: Duplicator plays 4 in $\mathcal{L}_7$
  Spoiler plays 6 in $\mathcal{L}_8$: Duplicator plays 6 in $\mathcal{L}_7$; spoiler cannot win
  Spoiler plays 7 in $\mathcal{L}_8$: Duplicator plays 6 in $\mathcal{L}_7$; spoiler cannot win
Other cases similar: Spoiler never wins

Theorem 11.10: The following are equivalent:

- $\mathcal{L}_m \sim_r \mathcal{L}_n$
- either (1) $m = n$, or (2) $m \geq 2^r - 1$ and $n \geq 2^r - 1$

Proof: For the $\implies$ direction, we show the contrapositive:

- If $m < n$ and $m < 2^r - 1$, then $\mathcal{L}_m \not\sim_r \mathcal{L}_n$
- If $m > n$ and $n < 2^r - 1$, then $\mathcal{L}_m \not\sim_r \mathcal{L}_n$

We define a winning strategy for the spoiler for the first case (the other is symmetrical).

Outline of the spoiler strategy:

- The spoiler always picks from the longer order $\mathcal{L}_n$
- The spoiler plays on an active interval of $\mathcal{L}_n$, initially the whole linear order
- In each step, the spoiler selects the middle point of the active interval
- Depending on the response by the duplicator, the spoiler updates the active interval to be the first or the second half of the previous interval

EF Games and Linear Orders

Proof (ctnd.): Detailed strategy of the spoiler:

- Initialise the spoiler's active interval (ℓ_n, u_n) with $\ell_n := 0$ and $u_n := n + 1$
- Initialise the duplicators's active interval (ℓ_m, u_m) with $\ell_m := 0$ and $u_m := m + 1$
- In each step:
    - The spoiler selects the central point $c := \ell_n + \lceil (u_n - \ell_n)/2 \rceil$
    - The duplicator answers by selecting some point d
    - If $d - \ell_m < u_m - d$, then we update $u_n := c$ and $u_m := d$; else
      if $d - \ell_m \geq u_m - d$, then we update $\ell_n := c$ and $\ell_m := d$

Analysis:

- To have a chance of winning, the duplicator must always select a d with $\ell_m < d < u_m$
  (not hard to see, since any non-initial value of ℓ_m and u_m comes from a previous selection round, and relative positions with respect to previously selected elements must be preserved)
- If an interval contains $2^j - 1$ elements, then selecting a center c and splitting the rest in half leads to a new active interval of size $(2^j - 2)/2 = 2^{j-1} - 1$
- Hence, after i rounds, the spoiler's active interval has size $\geq 2^{r-i} - 1$
- Similarly, the duplicator's active interval after i rounds is $< 2^{r-i} - 1$
- Therefore, after $r - 1$ rounds, the spoiler's active interval is $\geq 2^{r-r+1} - 1 = 1$, while the duplicator's is $< 2^{r-r+1} - 1 = 1$, i.e., empty – the duplicator looses

EF Games and Linear Orders

Theorem 11.10: The following are equivalent:

- $\mathcal{L}_m \sim_r \mathcal{L}_n$
- either (1) $m = n$, or (2) $m \geq 2^r - 1$ and $n \geq 2^r - 1$

Proof: It remains to show the $\Longleftarrow$ direction.

If $m = n$, Duplicator clearly wins, so assume $m \neq n$.

As before, it helps to think in terms of active intervals:

- For a round i ($0 \leq i \leq r$), we define potential interval boundaries on the two orders:

$$B_m(i) = \{a_1, \ldots, a_i\} \cup \{0, m + 1\} \qquad B_n(i) = \{b_1, \ldots, b_i\} \cup \{0, n + 1\}$$

- If Duplicator succeeds (as claimed), then we get a bijection $\iota : B_m(i) \to B_n(i)$ with $\iota(0) = 0$, $\iota(m + 1) = n + 1$, and $\iota(a_j) = b_j$ for all $1 \leq j \leq i$.
- In round $i + 1$, if Spoiler selects element $e \in \mathcal{L}_x$ ($x \in \{n, m\}$) that was not selected before, then e is in an active interval (ℓ_x, u_x) where $\ell_x = \max\{e' \in B_x(i) \mid e' < e\}$ and $u_x = \min\{e' \in B_x(i) \mid e' > e\}$.
- This interval (ℓ_x, u_x) corresponds to an interval $(\ell_y, u_y) = (\iota(\ell_x), \iota(u_x))$ on $\mathcal{L}_y$ ($y \in \{m, n\}, y \neq x$). Duplicator must select a suitable f in this interval (ℓ_y, u_y).

EF Games and Linear Orders

Proof (ctnd.): We show by induction that Duplicator can ensure the following:

> In any round i ($0 \leq i \leq r$), and for any interval (a, a') with $a, a' \in B_m(i)$ and corresponding interval $(b, b') = (\iota(a), \iota(a'))$ in $B_n(i)$, we have:
>
> (1) $a' - a = b' - b$, or
>
> (2) $a' - a \geq 2^{r-i}$ and $b' - b \geq 2^{r-i}$

Base: If $i = 0$, then $(a, a') = (0, m + 1)$ and the claim holds by the assumptions of the theorem.

Step: For round $i + 1$ ($0 < i + 1 \leq r$), assume the claim holds for i.

- If Spoiler picks a previously selected $e \in \mathcal{L}_x$, Duplicator picks the previously picked corresponding element in $\mathcal{L}_y$. The claim remains true.

- If Spoiler picks a new $e \in \mathcal{L}_x$ in active interval (ℓ_x, u_x), Duplicator selects an element f in the corresponding active interval (ℓ_y, u_y) in $\mathcal{L}_y$: ... (see next page)

EF Games and Linear Orders

Proof (ctnd.): We show by induction that Duplicator can ensure the following:

> In any round i ($0 \leq i \leq r$), and for any interval (a, a') with $a, a' \in B_m(i)$ and corresponding interval $(b, b') = (\iota(a), \iota(a'))$ in $B_n(i)$, we have:
>
> (1) $a' - a = b' - b$, or
>
> (2) $a' - a \geq 2^{r-i}$ and $b' - b \geq 2^{r-i}$

Step: For round $i + 1$ ($0 < i + 1 \leq r$), assume the claim holds for i.

- If Spoiler picks a new $e \in \mathcal{L}_x$ in active interval (ℓ_x, u_x), Duplicator selects an element f in the corresponding active interval (ℓ_y, u_y) in $\mathcal{L}_y$:

  - If $e - \ell_x < 2^{r-i-1}$ then $f = \ell_y + (e - \ell_x)$. The new intervals $(\ell_x, e)/(\ell_y, f)$ satisfy (1). Case A: if $(\ell_x, u_x)/(\ell_y, u_y)$ satisfies (1), then $(e, u_x)/(f, u_y)$ does too. Case B: if $(\ell_x, u_x)/(\ell_y, u_y)$ satisfies (2), then $(e, u_x)/(f, u_y)$ does too.

  - If $u_x - e < 2^{r-i-1}$ then $f = u_y - (u_x - e)$. Symmetric argument as in previous case. (Note: if both cases apply, then $(\ell_x, u_x)/(\ell_y, u_y)$ satisfies (1), and $b = \ell_y + (e - \ell_x) = \ell_y + (u_x - \ell_x) - (u_x - e) = \ell_y + (u_y - \ell_y) - (u_x - e) = u_y - (u_x - e)$.)

  - Else, $(\ell_x, u_x)/(\ell_y, u_y)$ satisfies (2), and Duplicator can pick any f with $\ell_y + 2^{r-i-1} \leq f \leq u_y - 2^{r-i-1}$. This is possible since, by (2), $u_y - \ell_y \geq 2^{r-i} = 2 * 2^{r-i-1}$. The new interval pairs $(\ell_x, e)/(\ell_y, f)$ and $(e, u_x)/(f, u_y)$ satisfy (2).

Theorem 11.10: The following are equivalent:

- $\mathcal{L}_m \sim_r \mathcal{L}_n$
- either (1) $m = n$, or (2) $m \geq 2^r - 1$ and $n \geq 2^r - 1$

Proof (ctnd.): We show by induction that Duplicator can ensure the following:

In any round i ($0 \leq i \leq r$), and for any interval (a, a') with $a, a' \in B_m(i)$ and corresponding interval $(b, b') = (\iota(a), \iota(a'))$ in $B_n(i)$, we have:

(1) $a' - a = b' - b$, or

(2) $a' - a \geq 2^{r-i}$ and $b' - b \geq 2^{r-i}$

In particular, our construction showed that Duplicator can maintain the correspondance of $B_m(i)$ and $B_n(i)$, and especially of $B_m(r)$ and $B_n(r)$.

In other words: Duplicator wins. □

FO-Definability of PARITY

Theorem 11.11: PARITY is not FO-definable for linear orders, hence it is not FO-definable for arbitrary databases.

Proof:

- Suppose for a contradiction that PARITY is FO-definable by some query φ.
- Let r be the quantifier rank of φ.
- Consider databases $\mathcal{L}_m$ and $\mathcal{L}_n$ with $m = 2^r$ and $n = 2^r + 1$.
- We know that $\mathcal{L}_m \sim_r \mathcal{L}_n$, and therefore $\mathcal{L}_m \equiv_r \mathcal{L}_n$.
- Hence, $\mathcal{L}_m \models \varphi$ if and only if $\mathcal{L}_n \models \varphi$.
- But $\mathcal{L}_m \in$ PARITY while $\mathcal{L}_n \notin$ PARITY.
- Therefore, φ does not FO-define PARITY. Contradiction. $\square$

Note: We might instead of linear orders have used cliques (of odd and even size). Linear orders are better to illustrate the exponential reach of quantifiers, but the EF game for cliques is simpler. Exercise: Show a suitable version of Theorem 10 for cliques (should become simpler), and use it to prove the above.

FO-Definability of CONNECTIVITY

The CONNECTIVITY problem over finite graphs is as follows:

CONNECTIVITY
- Input: A finite graph (relational structure with one binary relation "edge")
- Output: "true" if there is an (undirected) path between any pair of vertices

Theorem 11.12: CONNECTIVITY is not FO-definable.

Proof:

- Suppose for a contradiction that CONNECTIVITY is FO-definable using a query φ.
- We show that this would make PARITY FO-definable on linear orders.
- For a linear order $\mathcal{L}$ with order predicate $\leq$, we define a finite graph $\mathcal{G}(\mathcal{L})$ over a binary predicate "edge" such that $\mathcal{G}(\mathcal{L})$ is connected if and only if $\mathcal{L}$ has an odd number of elements.

Defining a Graph From a Linear Order

We use abbreviations for the following FO formulas:

$$\text{succ}[x, y] = (x \leq y) \wedge \neg(y \leq x) \wedge \qquad\qquad y \text{ is the successor of } x$$
$$\forall z.(z \leq x \vee y \leq z)$$
$$\text{min}[x] = \forall z.x \leq z \qquad\qquad x \text{ is the first element}$$
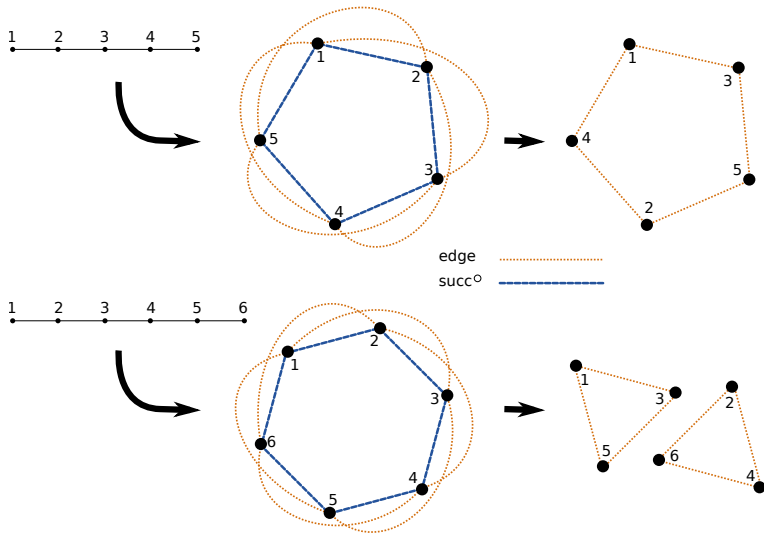$$\text{max}[x] = \forall z.z \leq x \qquad\qquad x \text{ is the last element}$$
$$\text{succ}^\circ[x, y] = \text{succ}[x, y] \vee (\text{max}[x] \wedge \text{min}[y]) \qquad \text{circular version of succ}$$
$$\text{edge}[x, y] = \exists z.(\text{succ}^\circ[x, z] \wedge \text{succ}^\circ[z, y]) \qquad \text{defining "edges" from the order}$$

Illustration: Graphs From Linear Orders

edge ....

succ° ─ ─ ─ ─ ─

Completing the Proof

Observation:

The graph $\mathcal{G}(\mathcal{L})$ is connected if and only if $\mathcal{L}$ has odd parity.

Finishing the proof:

Suppose for a contradiction that φ FO-defines CONNECTIVITY on graphs with predicate edge.

Then we can produce a formula φ' by replacing every atom of the form edge(s, t) in φ by the expression edge$[s, t]$ (a formula over $\leq$ only).

Then $\neg\varphi'$ FO-defines PARITY on linear orders.

Since PARITY is not FO-definable (not even on linear orders), no such φ can exist. $\qquad\square$

Beyond Linear Orders: Locality

Intuition: Duplicator can win an EF game if selected nodes have the same "neighbourhood"

$\rightsquigarrow$ let's define this for graphs (structures with binary predicates)

Definition 11.13: Consider a graph $\mathcal{G}$. For a natural number $d \geq 0$ and a vertex v, the d-neighbourhood of v, $N(v, d)$, is defined inductively:

- $N(v, 0) = \{v\}$
- $N(v, d + 1) = N(v, d) \cup \{w \mid w \text{ is a direct neighbour of some } w' \in N(v, d)\}$

Two vertices v and w have the same d-type if the subgraphs $\mathcal{G}|_{N(v,d)}$ and $\mathcal{G}|_{N(w,d)}$ are isomorphic.

Two graphs are d-equivalent if, for every d-type, they have the same number of d-neighbourhoods of this type.

Locality and FO-definability

A special case of Gaifman's Locality Theorem of first-order logic:

> **Theorem 11.14:** For every integer $r \geq 1$:
> - if $\mathcal{G}_1$ is 3^{r-1}-equivalent to $\mathcal{G}_2$
> - then $\mathcal{G}_1 \sim_r \mathcal{G}_2$, and thus $\mathcal{G}_1 \equiv_r \mathcal{G}_2$

$\rightsquigarrow$ Intuition: FO can only express local properties

How to show that a query mapping M can not be FO-defined:
- Let C_M be the class of all databases recognised by M
- Find sequences of graphs $\mathcal{I}_1, \mathcal{I}_2, \mathcal{I}_3, \ldots \in C_M$ and graphs $\mathcal{J}_1, \mathcal{J}_2, \mathcal{J}_3, \ldots \notin C_M$, such that $\mathcal{I}_i$ is i-equivalent to $\mathcal{J}_i$

$\rightsquigarrow$ for any formula φ (however large its quantifier rank r), there is a counterexample $\mathcal{I}_{3^{r-1}} \in C_M$ and $\mathcal{J}_{3^{r-1}} \notin C_M$ that φ cannot distinguish

CONNECTIVITY is not FO-definable (Proof 2)

Theorem 11.15: CONNECTIVITY is not FO-definable.

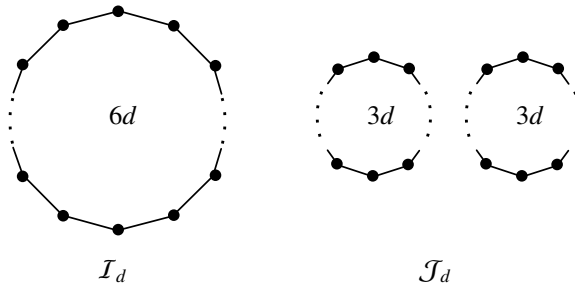**Proof:** counterexample for quantifier rank r: set $d = 3^r$
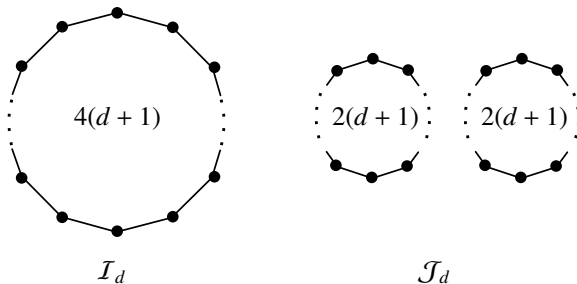
- the only d-type is a path of $2d + 1$ nodes
- $\mathcal{I}_d$ and $\mathcal{J}_d$ are d-equivalent □

2-COLOURABILITY

Theorem 11.16: 2-COLOURABILITY is not FO-definable.

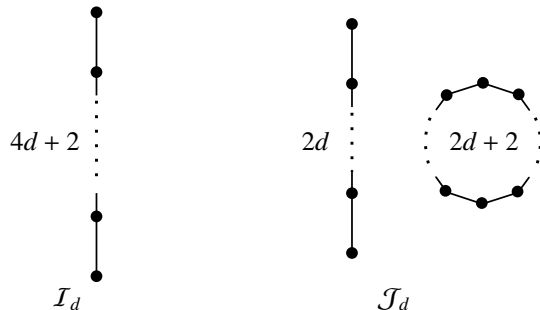**Proof:** counterexample for quantifier rank r: set $d = 3^r$ (odd number)

- the only d-type is a path of $2d + 1$ nodes
- $\mathcal{I}_d$ and $\mathcal{J}_d$ are d-equivalent   □

Theorem 11.17: ACYCLICITY is not FO-definable.

Proof: counterexample for quantifier rank r: set $d = 3^r$

- d-types are paths of $\leq 2d + 1$ nodes
- $\mathcal{I}_d$ and $\mathcal{J}_d$ are d-equivalent $\qquad\qquad$ □

Summary: Limits of FO-Queries

FO queries (and hence Relational Calculus) cannot express properties that require a "global" view:

- properties where one needs to follow paths
- properties where one needs to count elements

Remember Lecture 1?

> "Stops at distance 2 from Helmholtzstr."
>
> $$R_2 = \delta_{\text{To} \to \text{From}}(\pi_{\text{To}}(\text{Connect} \bowtie R_1))$$
>
> What about all stops reachable from Helmholtzstr.?

$\rightsquigarrow$ Not expressible in Relational Calculus

Yet, all examples we saw are in P
$\rightsquigarrow$ Is there another query language that could help us?

Summary and Outlook

FO-queries (and thus CQs) cannot express even all tractable query mappings ⤳ FO-definability

Showing that a query is not FO-definable requires some creativity ⤳ Ehrenfeucht-Fraïssé Games as one approach

FO-queries can only express "local" properties

Possible proof techniques:

- Ehrenfeucht-Fraïssé Games
- Locality Theorems
- For more approaches see
  Chapter 17 of [Abiteboul, Hull, Vianu 1994]

Open questions:
- If FO cannot express all tractable queries, what can?