

Complexity Theory
Exercise 5: Diagonalization

Exercise 5.1. Find the fault in the following proof of $P \neq NP$.

1. Assume that $P = NP$.
2. Then $SAT \in P$ and thus there exists a number $k \in \mathbb{N}$ such that $SAT \in DTime(n^k)$.
3. Because every language in NP is poly-time reducible to SAT , we have $NP \subseteq DTime(n^k)$.
4. It follows that $P \subseteq DTime(n^k)$.
5. By the Time Hierarchy Theorem there exist languages in $DTime(n^{k+1})$ that are not in $DTime(n^k)$, contradicting $P \subseteq DTime(n^k)$.
6. Therefore, $P \neq NP$.

Exercise 5.2. Show the following.

1. $TIME(2^n) = TIME(2^{n+1})$
2. $TIME_*(2^n) \subsetneq TIME_*(2^{2n})$
3. $NTIME(n) \subsetneq PSPACE$

Exercise 5.3. Define a function that is computable but not time-constructible.

Exercise 5.4. Consider the function $pad : \Sigma^* \times \mathbb{N} \rightarrow \Sigma^* \#^*$ defined as $pad(s, l) = s \#^j$, where $j = \max\{0, l - |s|\}$. For some language $A \subseteq \Sigma^*$ and function $f : \mathbb{N} \rightarrow \mathbb{N}$ define $pad(A, f) := \{pad(s, f(|s|)) \mid s \in A\}$.

Show all of the following statements.

1. If $A \in DTime(n^6)$, then $pad(A, n^2) \in DTime(n^3)$.
2. If $NEXPTIME \neq EXPTIME$, then $P \neq NP$.
3. For every $A \subseteq \Sigma^*$ and $k \in \mathbb{N}$, $A \in P$ iff $pad(A, n^k) \in P$.
4. $P \neq DSpace(n)$.
5. $NP \neq DSpace(n)$.

Exercise 5.5. You are given two oracles and one of them is the set **TQBF**, but you do not know which one. Design a polynomial algorithm that decides **TQBF** using these oracles.