

Advanced Topics in Complexity Theory

**Exercise 3: Function Problems**

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**Exercise 3.1** Show that FSAT is FNP-complete. For this first show that the function problem for

$$A_{\text{NPTM}} = \{ \langle M, w \rangle \mid M \text{ a polytime NTM accepting } w \}$$

is FNP-complete.

**Exercise 3.2** Consider the following function problem: given numbers  $a_1, \dots, a_n$  such that  $\sum_{i=1}^n a_i < 2^n - 1$ , find two different subsets  $S_1, S_2 \subseteq \{1, \dots, n\}$  such that

$$\sum_{i \in S_1} a_i = \sum_{i \in S_2} a_i.$$

Show that this problem is in TFNP.

**Exercise 3.3** Let  $G = (V, E)$  be an undirected graph with integer weights  $w$  on its edges. Think of the nodes as people, and of the edges as an indication of how much two people like each other (or not). A *state* of  $G$  is a mapping  $S: V \rightarrow \{+1, -1\}$ . We say that node  $i$  is *happy* in state  $S$  if

$$S(i) \cdot \sum_{\{i,j\} \in E} S(j)w(i,j) \geq 0.$$

The HAPPYNET problem is to find for each graph  $G$  a state in which each node is happy. Show that HAPPYNET is in TFNP. For this consider the mapping  $\varphi$  defined for states  $S$  by

$$\varphi(S) = \sum_{\{i,j\} \in E} S(i)S(j)w(i,j)$$

and suppose that some node  $i$  is unhappy. What happens to the value of  $\varphi(S)$  if one flips the current state of  $i$ ?

**Exercise 3.4** The goal of this exercise is to show that Primes is in NP (this is Pratt's Theorem). To this end we shall see that we can associate to every prime  $p$  a short certificate  $C(p)$  that can be checked in polynomial time and that no non-prime number has.

For this we make use of the following characterization of primes, proving which is not part of this exercise: a number  $p > 1$  is prime if and only if there is some  $1 \leq r < p$  such that  $r^{p-1} \equiv 1 \pmod{p}$  and  $r^{\frac{p-1}{q}} \not\equiv 1 \pmod{p}$  for all prime divisors  $q$  of  $p-1$ .

From this characterization, let us define the certificate  $C(p)$  for a prime  $p$  as

$$C(p) = (r, q_1, C(q_1), \dots, q_\ell, C(q_\ell))$$

where  $q_1, \dots, q_\ell$  are all prime divisors of  $p-1$ .

1. Verify that

$$C(67) = (2, 2, (1), 3, (2, 2, (1)), 11, (8, 2, (1), 5, (3, 2, (1))))).$$

2. Show that the length of  $C(p)$  is at most  $4(\log p)^2$  (i.e., polynomial in the length of  $p$ ).

3. Show that  $C(p)$  can be checked in polynomial time, i.e., it is checkable in polynomial time that

a)  $r^{p-1} \equiv 1 \pmod{p}$ ,

b)  $r^{\frac{p-1}{q}} \not\equiv 1 \pmod{p}$  for all  $q \in \{q_1, \dots, q_\ell\}$ ,

c) all  $q_1, \dots, q_\ell$  are prime, and

d)  $q_1, \dots, q_\ell$  are all prime factors of  $p - 1$ .