

Complexity Theory

Exercise 6: Diagonalisation and Alternation

Exercise 6.1. Show that Cook-reducibility is transitive. In other words, show that if **A** is Cook-reducible to **B** and **B** is Cook-reducible to **C**, then **A** is Cook-reducible to **C**.

Exercise 6.2. Show that there exists an oracle **C** such that $\text{NP}^{\text{C}} \neq \text{coNP}^{\text{C}}$.

Hint:

BAKKEI-CIJJ-2OJ0VASY TP6OIREEM IOI COIIB INSTEAD OF B.

What kind of Turing machines exist for languages in coNP^{C} ? Use the answer to adapt the proof of the

Exercise 6.3. Describe a polynomial-time ATM solving **EXACT INDEPENDENT SET**:

Input: Given a graph G and some number k .

Question: Does there exist a maximal independent set in G of size exactly k ?

Exercise 6.4. Consider the Japanese game *go-moku* that is played by two players **X** and **O** on a 19×19 board. Players alternately place markers on the board, and the first one to have five of its markers in a row, column, or diagonal wins.

Consider the generalised version of go-moku on an $n \times n$ board. Say that a *position* of go-moku is a placement of markers on such a board as it could occur during the game. Define

$$\mathbf{GM} = \{ \langle B \rangle \mid B \text{ is a position of go-moku where X has a winning strategy} \}.$$

Describe a polynomial-time ATM solving **GM**.

Exercise 6.5. Show that $\text{AEXP TIME} = \text{EXPSPACE}$.

* **Exercise 6.6.** Show that $\Sigma_2\text{QBF}$ is complete for $\Sigma_2\text{P}$. Generalise your argument to show that $\Sigma_i\text{QBF}$ is complete for $\Sigma_i\text{P}$ for all $i \geq 1$.

Exercise 6.7. Show that if $\text{P} = \text{NP}$, then $\text{P} = \text{PH}$.

Exercise 6.8. Describe a polynomial-time ATM solving **EXACT INDEPENDENT SET**.

$$\mathbf{EXACTIS} = \{ (G, k) \mid |S| = k \text{ for some independent set } S \text{ in } G \text{ and} \\ |S'| \leq k \text{ for every independent set } S' \text{ in } G \}$$

Find a level of the polynomial hierarchy where this problem is contained in.

Exercise 6.9. Consider the Japanese game *go-moku* that is played by two players X and O on a 19×19 board. Players alternately place markers on the board, and the first one to have five of its markers in a row, column, or diagonal wins.

Consider the generalised version of *go-moku* on an $n \times n$ board. Say that a *position* of *go-moku* is a placement of markers on such a board as it could occur during the game. Define

$$\mathbf{GM} = \{ \langle B \rangle \mid B \text{ is a position of go-moku where } X \text{ has a winning strategy} \}.$$

Describe a polynomial-time ATM solving \mathbf{GM} and informally argue why this problem is not in any level of the polynomial hierarchy.

Exercise 6.10. Show $\text{NP}^{\text{SAT}} \subseteq \Sigma_2\text{P}$.

Exercise 6.11. Show the following result: *If there is any k such that $\Sigma_k^{\text{P}} = \Sigma_{k+1}^{\text{P}}$ then $\Sigma_j^{\text{P}} = \Pi_j^{\text{P}} = \Sigma_k^{\text{P}}$ for all $j > k$, and therefore $\text{PH} = \Sigma_k^{\text{P}}$.*

Exercise 6.12. Show that $\text{PH} \subseteq \text{PSPACE}$.

Exercise 6.13. Let \mathbf{A} be a language and let \mathbf{F} be a finite set with $\mathbf{A} \cap \mathbf{F} = \emptyset$. Show that $\text{P}^{\mathbf{A}} = \text{P}^{\mathbf{A} \cup \mathbf{F}}$ and $\text{NP}^{\mathbf{A}} = \text{NP}^{\mathbf{A} \cup \mathbf{F}}$.