

Finite and Algorithmic Model Theory

Lecture 2 (Dresden 19.10.22, Short version with errors?)

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Today's agenda

Proviso: For simplicity we focus on finite, simple, undirected graphs today!

1. Asymptotic **probabilities** with examples.

Different perspective: What percentage of graphs verify a given FO sentence?

2. **Zero-One Law** of FO = **Probability** that a random structure satisfies φ is **always 0 or 1**.

3. Proof of the Zero-One Law for FO, based on Grädel's notes [HERE].

- Atomic k -types and **extensions axioms**. Theory \mathbb{EA} of extension axioms.
- Each extension axiom is **almost surely true**.
- \mathbb{EA} is **ω -categorical**, i.e. has exactly one countable model up to \cong , the Rado graph (**the random graph**).
- \mathbb{EA} is **complete**, i.e. for all $\varphi \in \text{FO}$ we have $\mathbb{EA} \models \varphi$ or $\mathbb{EA} \models \neg\varphi$.



Feel free to ask questions and interrupt me!

Don't be shy! If needed send me an email (bartosz.bednarczyk@cs.uni.wroc.pl) or approach me after the lecture!

Reminder: this is an advanced lecture. Target: people that had fun learning logic during BSc studies!

Asymptotic Probabilities

Proviso: For simplicity we focus on finite, simple, undirected graphs today!



We consider **random graphs**, according to the **uniform distribution**, i.e. every edge has probability $\frac{1}{2}$.

Let \mathcal{G}_n be the class of simple undirected graphs with n nodes. Of course $|\mathcal{G}_n| = 2^{\frac{n(n-1)}{2}}$.

Let \mathcal{P} be a property of graphs. Let $\mu_n(\mathcal{P}) =$ “probability that \mathcal{P} holds in a random graph with n nodes”.

$$\mu_n(\mathcal{P}) := \frac{|\{\mathcal{G} \in \mathcal{G}_n : \mathcal{G} \models \mathcal{P}\}|}{|\mathcal{G}_n|}$$

Asymptotic probability

$$\mu_\infty(\mathcal{P}) := \lim_{n \rightarrow \infty} \mu_n(\mathcal{P})$$

Examples

1. Take $\mathcal{P} :=$ “the graph is complete”. Then $\mu_\infty(\mathcal{P}) = \lim_{n \rightarrow \infty} \frac{|\{\mathcal{G} \in \mathcal{G}_n : \mathcal{G} \models \mathcal{P}\}|}{|\mathcal{G}_n|} = \lim_{n \rightarrow \infty} \frac{1}{2^{\frac{n(n-1)}{2}}} = 0$.

2. Take $\mathcal{P} :=$ “the graph has a triangle”. $\mu_3(\mathcal{P}) = \frac{1}{8}$. Since $\mu_{3n}(\mathcal{P}) \geq 1 - (1 - \frac{1}{8})^n$, we get $\mu_\infty(\mathcal{P}) = 1$.

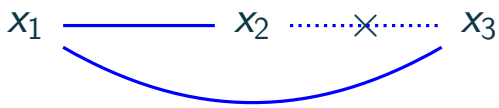
3. Take $\mathcal{P} :=$ “the graph has even number of edges”. $\mu_\infty(\mathcal{P}) = \frac{1}{2}$. Why?

$$\mu_\infty(\mathcal{P}) = \frac{|\{\mathcal{G} \in \mathcal{G}_n : \mathcal{G} \models \mathcal{P}\}|}{2^{\frac{n(n-1)}{2}}} = \frac{\sum_{i \geq 0} \binom{\frac{n(n-1)}{2}}{2i}}{2^{\frac{n(n-1)}{2}}} = [\text{Sum of Even Index Binomial Coeff.}] = \frac{2^{\frac{n(n-1)}{2} - 1}}{2^{\frac{n(n-1)}{2}}} = \frac{1}{2}$$

4. Take $\mathcal{P} :=$ “the graph has even number of nodes”. Then $\mu_\infty(\mathcal{P})$ does not exist.

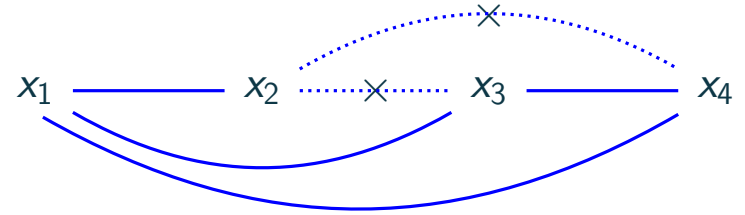
k-Types and Extension Axioms

A **k-type** is a conjunction of formulae with variables x_1, \dots, x_k such that for all $i \neq j$ we have $x_i \neq x_j$ and precisely one of $E(x_i, x_j)$ or $\neg E(x_i, x_j)$ as a conjunct.



$$s := x_1 \neq x_2 \wedge x_1 \neq x_3 \wedge x_2 \neq x_3 \wedge \\ E(x_1, x_2) \wedge \neg E(x_2, x_3) \wedge E(x_1, x_3)$$

⊆



$$t := x_1 \neq x_2 \wedge x_1 \neq x_3 \wedge x_2 \neq x_3 \wedge E(x_1, x_2) \wedge \neg E(x_2, x_3) \wedge E(x_1, x_3) \\ \wedge x_1 \neq x_4 \wedge x_2 \neq x_4 \wedge x_3 \neq x_4 \wedge E(x_1, x_4) \wedge \neg E(x_2, x_4) \wedge E(x_3, x_4)$$

A $(k + 1)$ -type t **extends** a k -type s if $\text{conjuncts}(s) \subseteq \text{conjuncts}(t)$ (c.f. the above picture).

An (s, t) -**extension axiom** $\sigma_{s,t}$ is $\forall x_1 \dots \forall x_k s(x_1, \dots, x_k) \rightarrow \exists x_{k+1} t(x_1, \dots, x_k, x_{k+1})$.

$$\mathbb{EA} := \left\{ \forall x \neg E(x, x), \forall xy E(x, y) \rightarrow E(y, x), \sigma_{s,t} \mid s \text{ is } k\text{-type, } t \text{ is } (k+1)\text{-type, } t \text{ extends } s \right\}$$

Why the theory \mathbb{EA} is important? Zero-One Law for $\text{FO}\{\{E\}\}$.

1. Every extension axiom $\sigma_{s,t}$ from \mathbb{EA} is **almost surely true**, i.e. $\mu_\infty(\sigma_{s,t}) = 1$ (Exercise).
2. By Compactness, it follows that $\mathbb{EA} \models \varphi$ implies $\mu_\infty(\varphi) = 1$ (TODO).
3. The theory \mathbb{EA} is **ω -categorical**, i.e. has exactly one countable model up to \cong (TODO).
4. Thus \mathbb{EA} is **complete**, i.e. for all $\varphi \in \text{FO}$ we have $\mathbb{EA} \models \varphi$ or $\mathbb{EA} \models \neg\varphi$ (TODO).

Theorem (Glebskii et al. 1969, Fagin 1976)

For every formula $\varphi \in \text{FO}\{\{E\}\}$ we have that $\mu_\infty(\varphi)$ is either 0 or 1.

Proof

Take any $\varphi \in \text{FO}\{\{E\}\}$. By (4) either $\mathbb{EA} \models \varphi$ or $\mathbb{EA} \models \neg\varphi$. If $\mathbb{EA} \models \varphi$ then by (2) we have $\mu_\infty(\varphi) = 1$. Otherwise $\mathbb{EA} \models \neg\varphi$, so by (2) we infer $\mu_\infty(\neg\varphi) = 1$, which leads to $\mu_\infty(\varphi) = 1 - \mu_\infty(\neg\varphi) = 0$. ■

Applications?

- Evenness of the number of nodes/edges not $\text{FO}\{\{E\}\}$ -definable.
- No information about connectivity because $\mu_\infty(\text{"graph is connected"}) = 0$.

Proof of $\mathbb{EA} \models \varphi$ implies $\mu_\infty(\varphi) = 1$ (assuming that $\forall \sigma \in \mathbb{EA} \mu_\infty(\sigma) = 1$).

Handy observations for all $\alpha, \beta, \gamma \in \text{FO}[\{E\}]$ and all $n \in \mathbb{N}$:

$$\mu_n(\neg\alpha) = 1 - \mu_n(\alpha)$$



Compactness: $\mathbb{EA} \models \varphi$ implies
there is $\mathbb{EA}_0 \subseteq_{\text{fin}} \mathbb{EA}$ implying φ



$$\mu_n(\beta \vee \gamma) \leq \mu_n(\beta) + \mu_n(\gamma).$$



Proof

Goal: To show $\mu_\infty(\varphi) = 1$ it suffices to show that $\mu_n(\neg\varphi) \rightarrow 0$ when $n \rightarrow \infty$.

Assume $\mathbb{EA} \models \varphi$. By compactness, there is a finite $\mathbb{EA}_0 \subseteq \mathbb{EA}$ such that $\mathbb{EA}_0 \models \varphi$.

So $\mu_n(\varphi) \geq \mu_n(\wedge \mathbb{EA}_0)$, thus $\mu_n(\neg \wedge \mathbb{EA}_0) \geq \mu_n(\neg\varphi)$.

Moreover (by our assumption), $\mu_n(\neg\sigma) = 1 - \mu_n(\sigma)$ tends to 0 when $n \rightarrow \infty$.

$$\mu_n(\neg\varphi) \leq \mu_n(\neg \wedge \mathbb{EA}_0) = \mu_n\left(\bigvee_{\sigma \in \mathbb{EA}_0} \neg\sigma\right) \leq \sum_{\sigma \in \mathbb{EA}_0} \mu_n(\neg\sigma)$$

The sum $\sum_{\sigma \in \mathbb{EA}_0} \mu_n(\neg\sigma)$ converges to 0 for $n \rightarrow \infty$, concluding $\mu_\infty(\varphi) = 1$.

\mathbb{EA} is satisfiable and complete (assuming ω -categoricity)

- Note that $\mathbb{EA} \not\models \forall x \perp$ (due to $\mu_\infty(\forall x \perp) = 0$). So \mathbb{EA} have a model (UnSAT theory entails everything).

\mathbb{EA} is complete (assuming ω -categoricity), i.e. for all φ we either have $\mathbb{EA} \models \varphi$ or $\mathbb{EA} \models \neg\varphi$.

Proof

Assume that \mathbb{EA} is not complete. Thus we have $\mathfrak{A} \models \varphi$ and $\mathfrak{B} \models \neg\varphi$ that are both models of \mathbb{EA} .

Since $|\mathbb{EA}| = \aleph_0$, by Löwenheim-Skolem we can assume w.l.o.g. that \mathfrak{A} and \mathfrak{B} are also countably-infinite.

But then, by ω -categoricity of \mathbb{EA} , we infer $\mathfrak{A} \cong \mathfrak{B}$.

Thus $\mathfrak{A} \models \varphi$ and $\mathfrak{A} \models \neg\varphi$ (since $\mathfrak{B} \models \neg\varphi$). A contradiction! ■

Ad absurdum



Löwenheim-Skolem



ω -categoricity



\cong preserves \models



Today's final boss: \mathbb{EA} is ω -categorical

\mathbb{EA} is ω -categorical, i.e. has precisely one countably-infinite model.

Take any two countably-inf models $\mathfrak{A}, \mathfrak{B}$ of \mathbb{EA} with the domains $A := \{a_1, a_2, \dots\}$ and $B := \{b_1, b_2, \dots\}$.

Goal: We will show that $\mathfrak{A} \cong \mathfrak{B}$ by producing an infinite growing **sequence of partial isomorphisms** p_0, p_1, \dots .

The union $\bigcup_{i=0}^{\infty} p_i$ will be the desired **isomorphism**. Start from $p_0 := \emptyset$.

Assume that a **partial isomorphism** $p_n = \{a_{i_1} \mapsto b_{i_1}, a_{i_2} \mapsto b_{i_2}, \dots, a_{i_n} \mapsto b_{i_n}\}$ is given. Goal: define p_{n+1} .

If $n+1$ is even, we will **select some element from \mathfrak{A}** (otherwise proceed analogously in \mathfrak{B} , proof omitted).

Take $a_k \in A$, for which k is the **smallest index** so that a_k does **not appear in p_n** . What do we know about \bar{a} ?

There are **unique n - and $(n+1)$ -types s and t** such that $s \subseteq t$, $\mathfrak{A} \models s(a_{i_1}, \dots, a_{i_n})$, and $\mathfrak{A} \models t(a_{i_1}, \dots, a_{i_n}, a_k)$.

Since p_n is a partial isomorphism, we have $\mathfrak{B} \models s(b_{i_1}, \dots, b_{i_n})$. But $\sigma_{s,t} \in \mathbb{EA}$ and $\mathfrak{B} \models \mathbb{EA}$!

Thus $\mathfrak{B} \models \sigma_{s,t} := \forall x_1 \dots \forall x_n s(x_1, \dots, x_n) \rightarrow \exists x_{n+1} t(x_1, \dots, x_n, x_{n+1})$.

So **there is an $b \in B$** so that $\mathfrak{B} \models t(b_{i_1}, \dots, b_{i_n}, b)$. Continue from $p_{n+1} := p_n \cup \{(a_k \mapsto b)\}$. ■

induction



back and forth



first not yet covered



exploit types realized by \bar{a}



ind. ass.



$\mathfrak{B} \models \mathbb{EA}$



Choose a witness



Extra: The Random Graph

We proved that \mathbb{EA} has a model unconstructively.

Can we describe the countable model of \mathbb{EA} ?

Let $\mathfrak{G} = (V, E)$ be a graph such that $V = \mathbb{N}_+$ and $(i, j) \in E^{\mathfrak{G}}$ iff $p_i \mid j$ or $p_j \mid i$ (p_i is the i -th prime number)

Lemma

$$\mathfrak{G} \models \sigma_{s,t} := \forall x_1 \dots \forall x_n s(x_1, \dots, x_n) \rightarrow \exists x_{n+1} t(x_1, \dots, x_n, x_{n+1})$$

Proof

Take any a_1, \dots, a_k such that $\mathfrak{G} \models s(a_1, \dots, a_k)$. Goal: Find a_{k+1} such that $\mathfrak{G} \models t(a_1, \dots, a_k, a_{k+1})$.

We divide indices $1, 2, \dots, k$ into $\text{Con} := \{i \mid E(x_i, x_{k+1})\}$ and $\text{DisC} := \{i \mid \neg E(x_i, x_{k+1})\}$.

Thus, our a_{k+1} must be connected to all a_i with $i \in \text{Con}$ and disconnected from all a_i with $i \in \text{DisC}$.

$$a_{k+1} := \prod_{i \in \text{Con}} p_{a_i} \cdot q, \text{ where } q \text{ is any prime number bigger than } \prod_{i=1}^k p_{a_i}$$

And now it is easy to check our choice of a_{k+1} is correct. ■

Divide x_1, x_2, \dots, x_k biased on type connections with $k+1$ (Dis)connected with $x \approx$ (non)dividable by the x -th prime number



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