

Weak Completion Semantics and Abduction

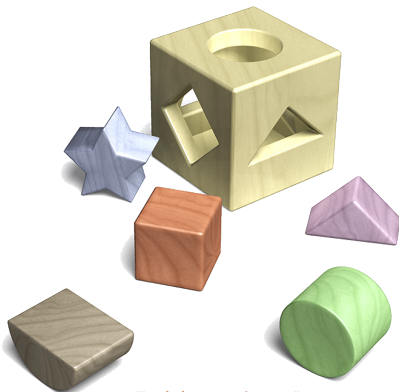
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- ▶ The Suppression Task
- ▶ Rules, Facts and Assumptions
- ▶ Weak Completion Semantics
- ▶ Abduction



"Logic is everywhere ..."



The Suppression Task

- Byrne: Suppressing Valid Inferences with Conditionals. Cognition 31, 61-83: 1989

- C_1 If she has an essay to write, then she will go to the library
 C_2 if she has a textbook to read, then she will go to the library
 C_3 If the library is open, then she will go to the library
 E She has an essay to write
 L She will go to the library

	E	$\neg E$	L	$\neg L$
C_1	96%	46%	71%	92%
$C_1 \& C_2$	96%	4%	13%	96%
$C_1 \& C_3$	38%	63%	54%	33%
	(L)	($\neg L$)	(E)	($\neg E$)



The Suppression Task – Modus Ponens

- ▶ Stenning, van Lambalgen: Human Reasoning and Cognitive Science. MIT Press: 2008

▶ Programs

$$\begin{aligned} e &\leftarrow \top \\ \ell &\leftarrow e \wedge \neg ab_1 \\ ab_1 &\leftarrow \perp \end{aligned}$$

fact
rule
assumption

definition of e
definition of ℓ
 ab_1 is assumed to be false

▶ Weakly completed programs & least models

$$\begin{aligned} e &\leftrightarrow \top \\ \ell &\leftrightarrow e \wedge \neg ab_1 \\ ab_1 &\leftrightarrow \perp \end{aligned}$$

<i>true</i>	<i>false</i>
e	ab_1
ℓ	

$\Phi \uparrow 1$
 $\Phi \uparrow 2$

▶ Computing logical consequences with respect to least models

$$\mathcal{M} = \langle \{e, \ell\}, \{ab_1\} \rangle \models_{wcs} \ell$$

- ▶ Łukasiewicz: O logice trójwartościowej. Ruch Filozoficzny 5, 169-171: 1920
- H., Kencana Ramli: Logic Programs under Three-Valued Łukasiewicz's Semantics LNCS 5649, 464-478: 2009



The Suppression Task – Alternative Argument

► Programs

$e \leftarrow \top$	fact	definition of e
$\ell \leftarrow e \wedge \neg ab_1$	rule	definition of ℓ
$ab_1 \leftarrow \perp$	assumption	ab_1 is assumed to be false
$\ell \leftarrow t \wedge \neg ab_2$	rule	definition of ℓ
$ab_2 \leftarrow \perp$	assumption	ab_2 is assumed to be false

► Weakly completed programs & least models

$e \leftrightarrow \top$	true	false	
$\ell \leftrightarrow (e \wedge \neg ab_1) \vee (t \wedge \neg ab_2)$	e	ab_1	
$ab_1 \leftrightarrow \perp$		ab_2	$\Phi \uparrow 1$
$ab_2 \leftrightarrow \perp$	ℓ		$\Phi \uparrow 2$

► Computing logical consequences with respect to least models

$$\mathcal{M} = \langle \{e, \ell\}, \{ab_1, ab_2\} \rangle \models_{wcs} \ell$$



The Suppression Task – Additional Argument

► Programs

$e \leftarrow \top$	fact	definition of e
$\ell \leftarrow e \wedge \neg ab_1$	rule	definition of ℓ
$ab_1 \leftarrow \perp$	assumption	ab_1 is assumed to be false
$\ell \leftarrow o \wedge \neg ab_3$	rule	definition of ℓ
$ab_3 \leftarrow \perp$	assumption	ab_3 is assumed to be false
$ab_1 \leftarrow \neg o$	rule	definition of ab_1
$ab_3 \leftarrow \neg e$	rule	definition of ab_3

► Weakly completed programs & least models

$e \leftrightarrow \top$	<i>true</i>	<i>false</i>	
$\ell \leftrightarrow (e \wedge \neg ab_1) \vee (o \wedge \neg ab_3)$	<u>e</u>		$\Phi \uparrow 1$
$ab_1 \leftrightarrow \perp \vee \neg o$		<u>ab_3</u>	$\Phi \uparrow 2$
$ab_3 \leftrightarrow \perp \vee \neg e$			

► Computing logical consequences with respect to least models

$$\mathcal{M} = \langle \{e\}, \{ab_3\} \rangle \not\models_{wcs} \ell$$



The Suppression Task – Denial of the Antecedent

► Programs

$$\begin{aligned} e &\leftarrow \perp \\ \ell &\leftarrow e \wedge \neg ab_1 \\ ab_1 &\leftarrow \perp \end{aligned}$$

► Weakly completed programs & least models

$$\begin{aligned} e &\leftrightarrow \perp \\ \ell &\leftrightarrow e \wedge \neg ab_1 \\ ab_1 &\leftrightarrow \perp \end{aligned}$$

<i>true</i>	<i>false</i>
	e
	ab_1
	ℓ

► Computing logical consequences with respect to least models

$$\mathcal{M} = \langle \emptyset, \{e, ab_1, \ell\} \rangle \models_{wcs} \neg \ell$$



The Suppression Task – Affirmation of the Consequent

► Programs

$$\begin{aligned} \ell &\leftarrow \top \\ \ell &\leftarrow e \wedge \neg ab_1 \\ ab_1 &\leftarrow \perp \end{aligned}$$

► Weakly completed programs & least models

$$\begin{aligned} \ell &\leftrightarrow \top \vee (e \wedge \neg ab_1) \\ ab_1 &\leftrightarrow \perp \end{aligned}$$

<i>true</i>	<i>false</i>
ℓ	ab_1

► Computing logical consequences with respect to least models

$$\mathcal{M} = \langle \{\ell\}, \{ab_1\} \rangle \not\models_{wcs} e$$

► Byrne 1989 **most humans conclude e !**



Abduction

► Programs & observations

$$\begin{array}{l} \ell \leftarrow e \wedge \neg ab_1 \\ ab_1 \leftarrow \perp \end{array} \qquad \ell$$

► Abducibles

$$e \leftarrow \top \qquad e \leftarrow \perp$$

► Weakly completed programs plus explanations & least models

$$\begin{array}{l} \ell \leftrightarrow e \wedge \neg ab_1 \\ ab_1 \leftrightarrow \perp \\ e \leftrightarrow \top \end{array} \qquad \begin{array}{cc} \textit{true} & \textit{false} \\ \hline e & ab_1 \\ \hline \ell & \end{array}$$

► Computing logical consequences with respect to least models

$$\mathcal{M} = \langle \{e, \ell\}, \{ab_1\} \rangle \models_{wcs} e$$

- H., Philipp, Wernhard: An Abductive Model for Human Reasoning. In: Proc. Tenth Int. Symposium on Logical Formalizations of Commonsense Reasoning: 2011



Alternative Arguments and Affirmation of the Consequent

► Programs & observations

$$\begin{array}{ll}
 \ell \leftarrow e \wedge \neg ab_1 & \ell \\
 ab_1 \leftarrow \perp & \\
 \ell \leftarrow t \wedge \neg ab_2 & \\
 ab_2 \leftarrow \perp &
 \end{array}$$

► Abducibles

$$e \leftarrow \top \quad t \leftarrow \top \quad e \leftarrow \perp \quad t \leftarrow \perp$$

► Weakly completed programs plus explanations & least models

$$\begin{array}{ll}
 \ell \leftrightarrow (e \wedge \neg ab_1) \vee (t \wedge \neg ab_2) & \begin{array}{cc} \text{true} & \text{false} \\ \hline e & ab_1 \\ & ab_2 \\ \hline \ell & \end{array} & \begin{array}{cc} \text{true} & \text{false} \\ \hline t & ab_1 \\ & ab_2 \\ \hline \ell & \end{array} \\
 ab_1 \leftrightarrow \perp & & \\
 ab_2 \leftrightarrow \perp & & \\
 e \leftrightarrow \top \quad \text{or} \quad t \leftrightarrow \top & &
 \end{array}$$

► Computing skeptical consequences with respect to both models

▷ **e does not follow**

- Dietz, H., Ragni: A Computational Logic Approach to the Suppression Task
Proc. COGSCI, 1500-1505: 2012



Usually Birds Fly

► Programs & observations

$$\begin{array}{ll}
 \text{fly}(X) & \leftarrow \text{bird}(X) \wedge \neg \text{ab}(X) & \neg \text{fly}(\text{tweety}) \\
 \text{ab}(X) & \leftarrow \perp \\
 \text{bird}(\text{tweety}) & \leftarrow \top \\
 \text{bird}(\text{jerry}) & \leftarrow \top
 \end{array}$$

► Abducibles

 \emptyset

► Weakly completed programs plus explanations & least models

$$\begin{array}{ll}
 \text{fly}(\text{tweety}) & \leftrightarrow \text{bird}(\text{tweety}) \wedge \neg \text{ab}(\text{tweety}) \\
 \text{fly}(\text{jerry}) & \leftrightarrow \text{bird}(\text{jerry}) \wedge \neg \text{ab}(\text{jerry}) \\
 \text{ab}(\text{tweety}) & \leftrightarrow \perp \\
 \text{ab}(\text{jerry}) & \leftrightarrow \perp \\
 \text{bird}(\text{tweety}) & \leftrightarrow \top \\
 \text{bird}(\text{jerry}) & \leftrightarrow \top
 \end{array}$$

<i>true</i>	<i>false</i>
<i>bird(tweety)</i>	<i>ab(tweety)</i>
<i>bird(jerry)</i>	<i>ab(jerry)</i>
<i>fly(tweety)</i>	
<i>fly(jerry)</i>	

► We cannot explain the observation!



Abducibles Revised

► Programs & observations

$$\begin{array}{ll}
 \text{fly}(X) \leftarrow \text{bird}(X) \wedge \neg \text{ab}(X) & \neg \text{fly}(\text{tweety}) \\
 \text{ab}(X) \leftarrow \perp & \\
 \text{bird}(\text{tweety}) \leftarrow \top & \\
 \text{bird}(\text{jerry}) \leftarrow \top &
 \end{array}$$

► Abducibles

$$\text{ab}(\text{tweety}) \leftarrow \top \qquad \text{ab}(\text{jerry}) \leftarrow \top$$

► Weakly completed programs plus explanations & least models

$ \begin{array}{ll} \text{fly}(\text{tweety}) \leftrightarrow \text{bird}(\text{tweety}) \wedge \neg \text{ab}(\text{tweety}) \\ \text{fly}(\text{jerry}) \leftrightarrow \text{bird}(\text{jerry}) \wedge \neg \text{ab}(\text{jerry}) \\ \text{ab}(\text{tweety}) \leftrightarrow \perp \vee \top \\ \text{ab}(\text{jerry}) \leftrightarrow \perp \\ \text{bird}(\text{tweety}) \leftrightarrow \top \\ \text{bird}(\text{jerry}) \leftrightarrow \top \end{array} $	<table border="1"> <thead> <tr> <th><i>true</i></th> <th><i>false</i></th> </tr> </thead> <tbody> <tr> <td><i>bird(tweety)</i></td> <td><i>ab(jerry)</i></td> </tr> <tr> <td><i>bird(jerry)</i></td> <td></td> </tr> <tr> <td><i>ab(tweety)</i></td> <td></td> </tr> <tr> <td><i>fly(jerry)</i></td> <td><i>fly(tweety)</i></td> </tr> </tbody> </table>	<i>true</i>	<i>false</i>	<i>bird(tweety)</i>	<i>ab(jerry)</i>	<i>bird(jerry)</i>		<i>ab(tweety)</i>		<i>fly(jerry)</i>	<i>fly(tweety)</i>
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<i>ab(tweety)</i>											
<i>fly(jerry)</i>	<i>fly(tweety)</i>										

► Dietz Saldanha, H., Pereira

Contextual Reasoning: Usually Birds can Abductively Fly. TUD: 2017 (submitted)



Summary

- ▶ Usually,
 - ▷ \mathcal{P} denotes a program, i.e. a finite set of facts, rules and assumptions
 - ▷ $\mathbf{g}\mathcal{P}$ denotes the ground instance of \mathcal{P}
 - ▷ $\mathbf{wc}\mathcal{P}$ denotes the weak completion of $\mathbf{g}\mathcal{P}$
 - ▷ $\mathcal{M}_{\mathcal{P}}$ denotes the least model of $\mathbf{wc}\mathcal{P}$
 - ▷ $\mathcal{A}_{\mathcal{P}}$ denotes the set of abducibles of $\mathbf{g}\mathcal{P}$
 - ▷ \mathcal{IC} denotes a finite set of integrity constraints of the form $\bigcup \leftarrow L_1 \wedge \dots \wedge L_n$
- ▶ **Weak Completion Semantics** is the approach
 - ▷ to consider weakly completed programs
 - ▷ to compute their least models
 - ▷ and to reason with respect to these models

