General Acyclicity and Cyclicity Notions for the Disjunctive Skolem Chase (Extended Abstract)

Lukas Gerlach¹ David Carral²

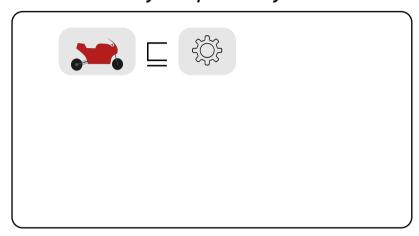
¹Knowledge-Based Systems, TU Dresden ²LIRMM, Inria, University of Montpellier, CNRS

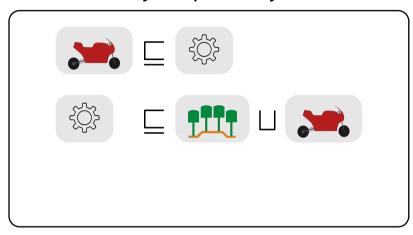
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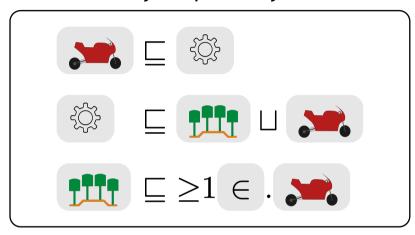


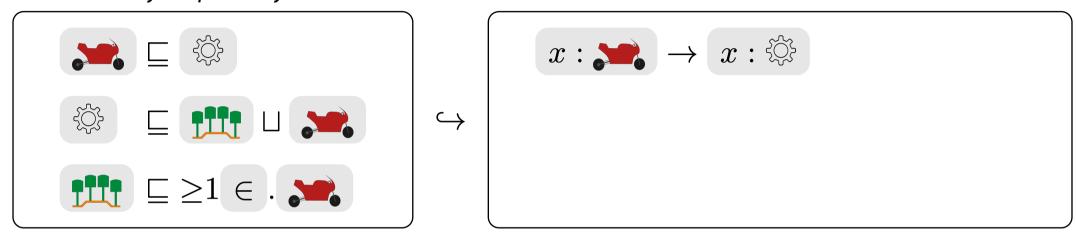


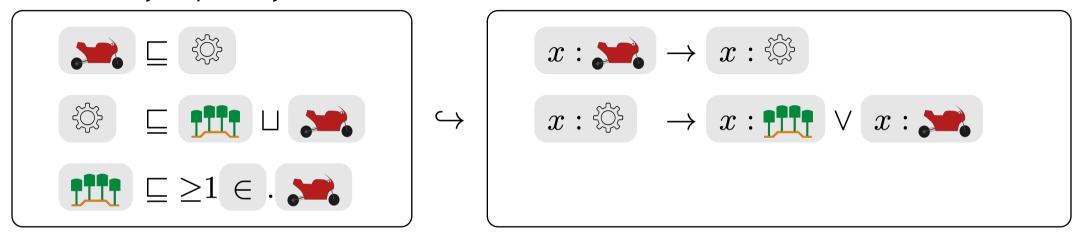


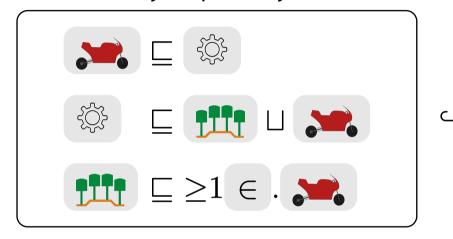


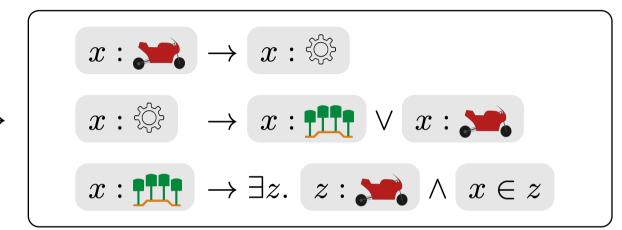


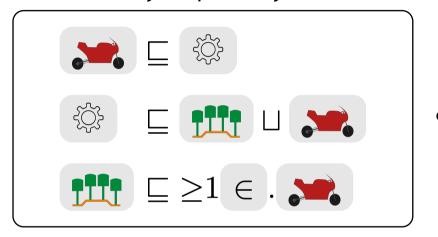


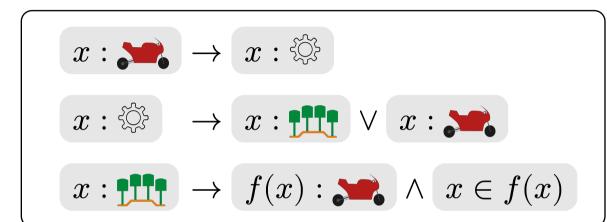


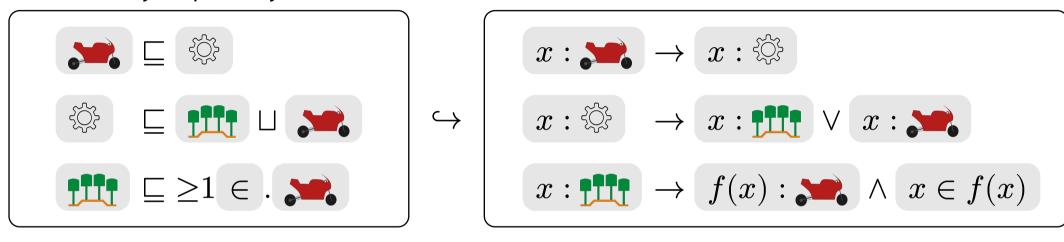


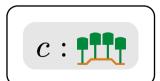


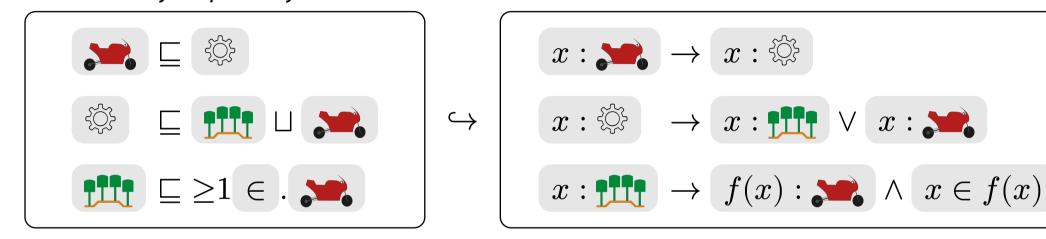


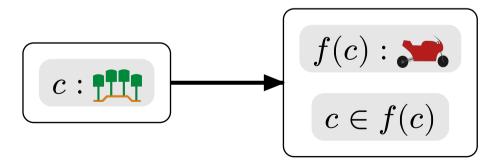


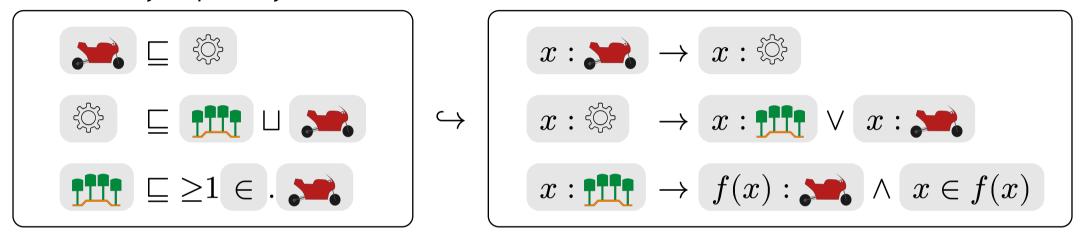


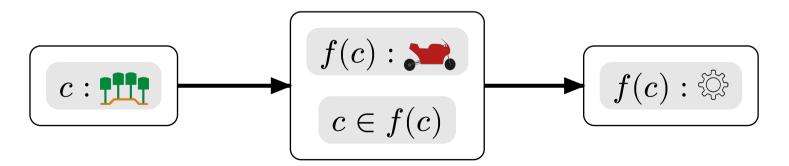


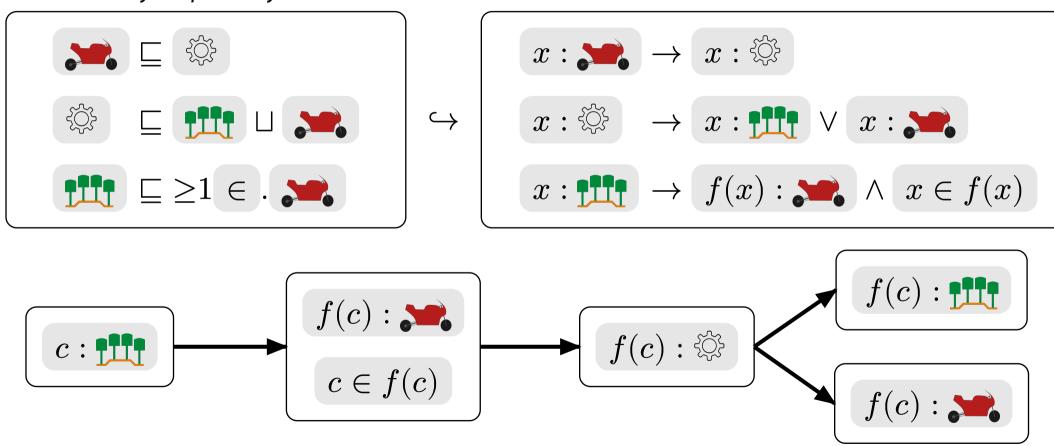


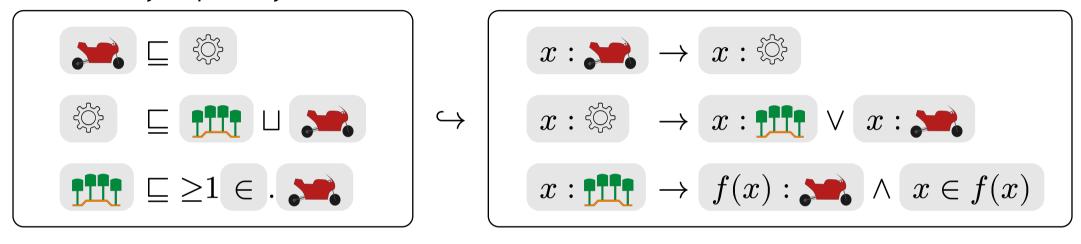


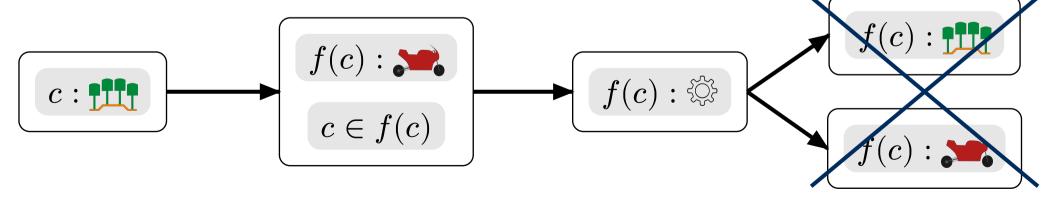


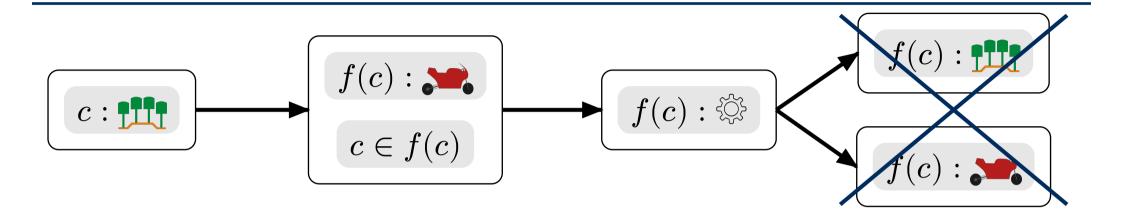


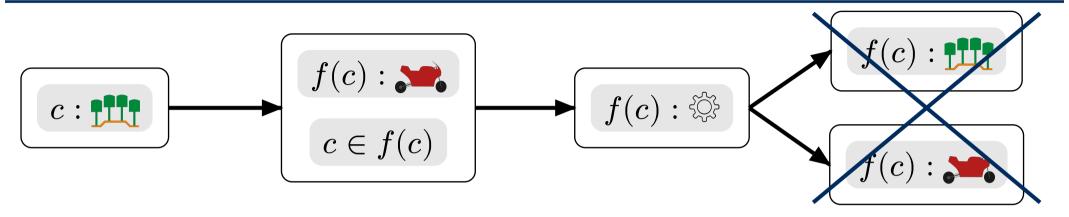




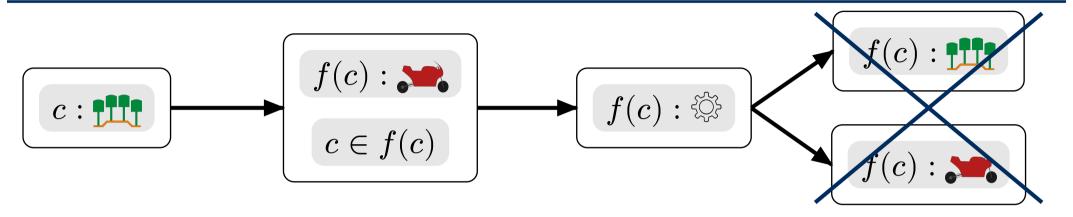




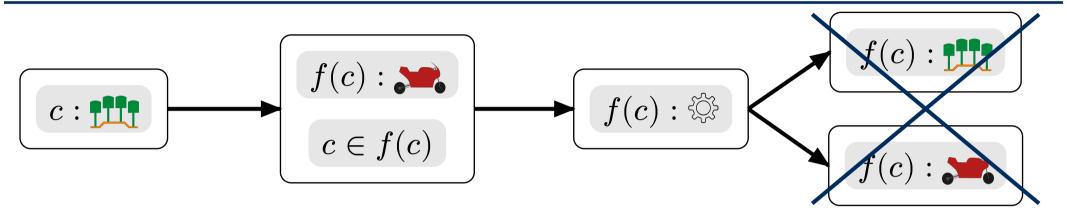




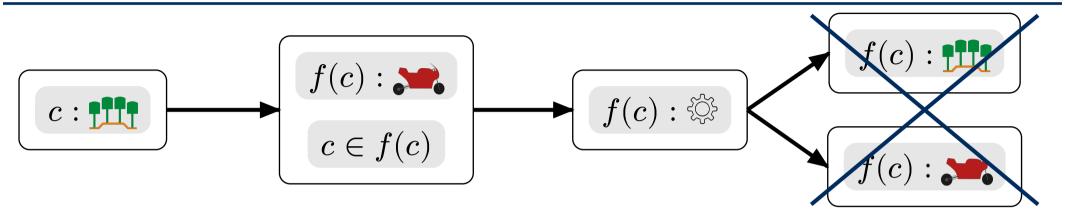
© Leaves of tree form *universal model set* (can answer queries) [Bou+16]



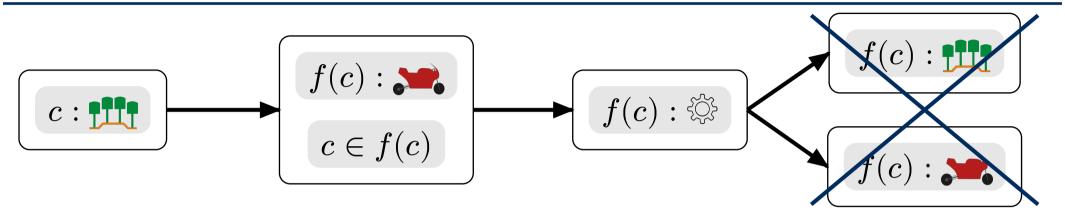
- Leaves of tree form universal model set (can answer queries) [Bou+16]
- But query answering/entailment is undecidable [BV81]



- Eaves of tree form *universal model set* (can answer queries) [Bou+16]
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- Chase termination is also undecidable [GM14, GO18]



- Example 2 Leaves of tree form *universal model set* (can answer queries) [Bou+16]
- But query answering/entailment is undecidable [BV81]
- Chase termination is also undecidable [GM14, G018]
- MFA/MFC can sometimes detect (non-)termination [Cue+13, CDK17]

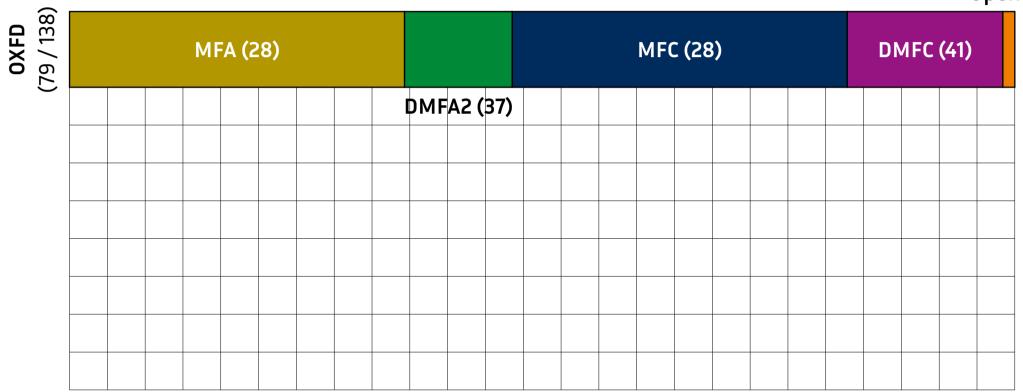


- Example 2 Leaves of tree form *universal model set* (can answer queries) [Bou+16]
- But query answering/entailment is undecidable [BV81]
- Chase termination is also undecidable [GM14, G018]
- MFA/MFC can sometimes detect (non-)termination [Cue+13, CDK17]
- We use ideas from RMFA/RMFC (for a different chase variant) [CDK17] to improve MFA/MFC for disjunctions: DMFA(2)/DMFC [GC23a]

Is it any good?

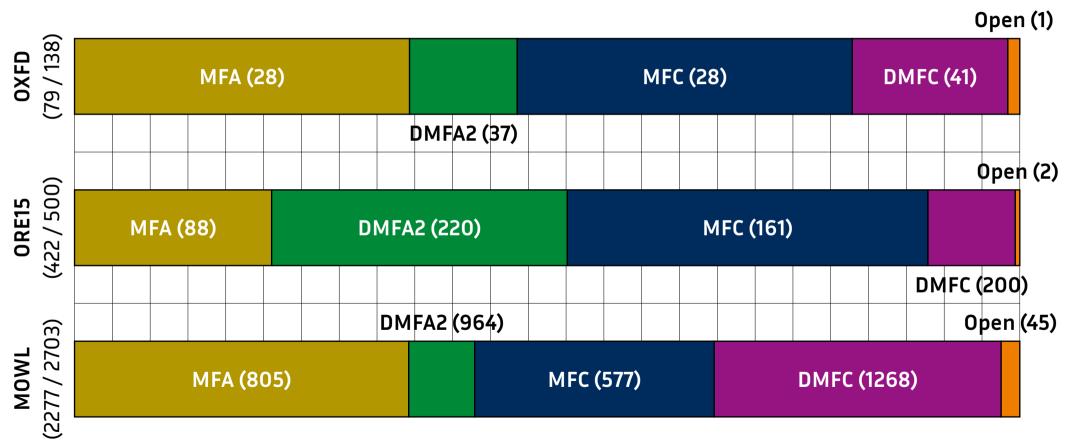
cat my-ruleset | docker run --rm -i registry.gitlab.com/m0nstr/dmfa-checker -t [non_]termination -cv skolem [-disj] [-depth 2]

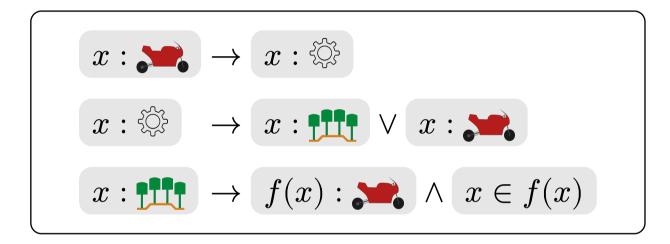
Open (1)



Is it any good?

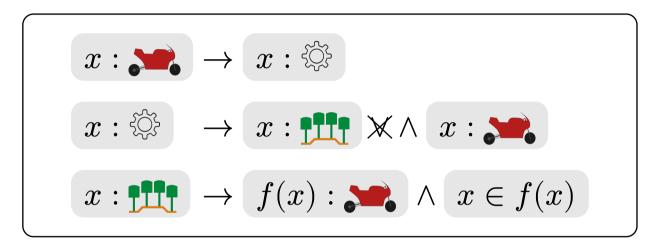
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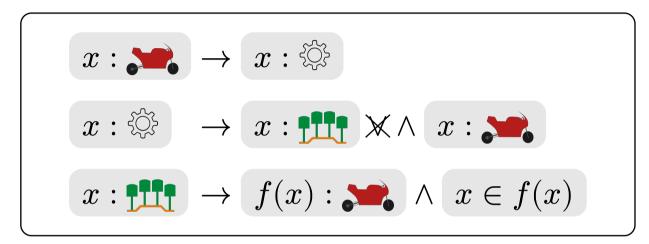


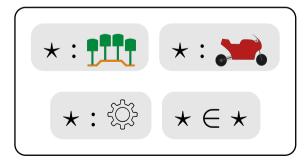
We want to know if the chase terminates on all databases for a rule set.

1. Replace Disjunctions

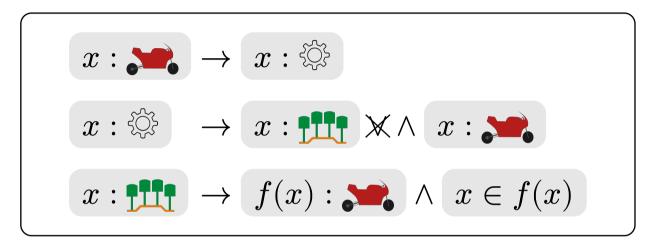


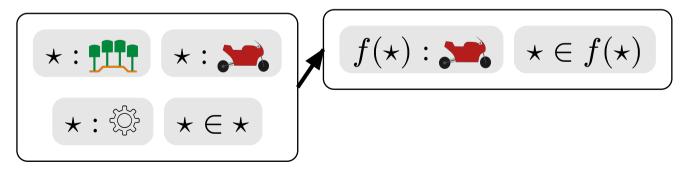
- 1. Replace Disjunctions
- 2. Use Critical Database



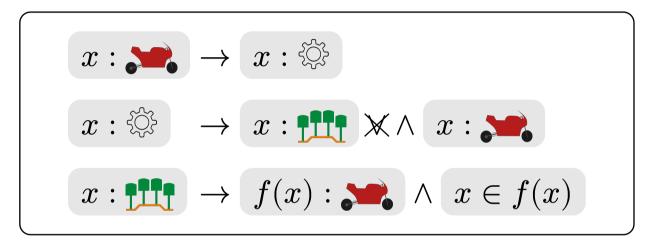


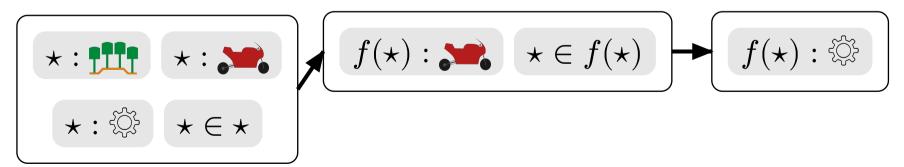
- 1. Replace Disjunctions
- 2. Use Critical Database
- 3. Run Chase



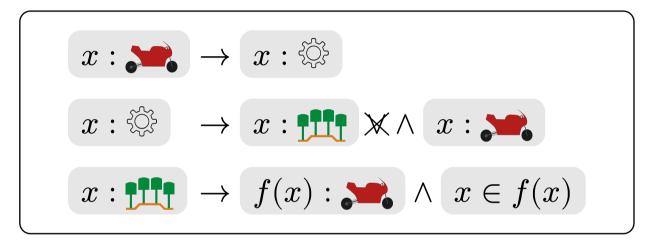


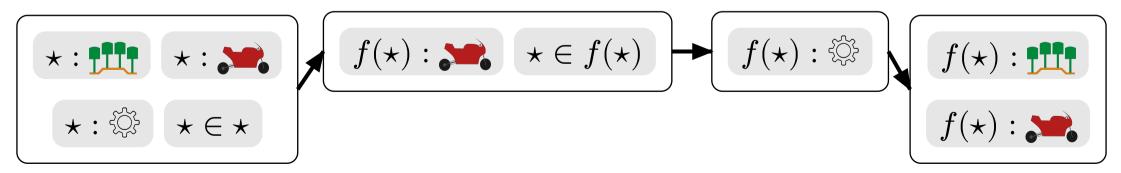
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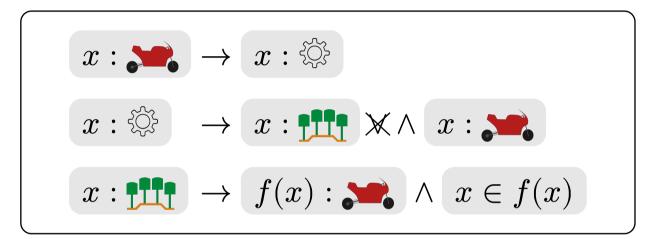


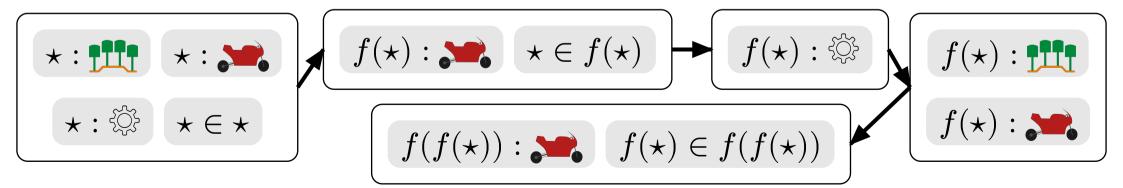
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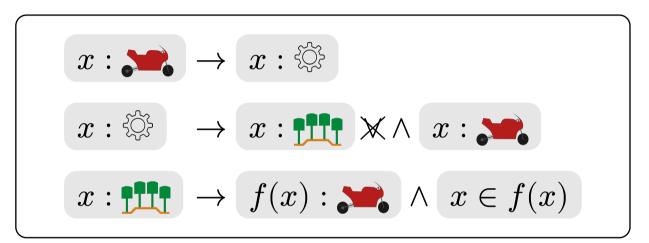


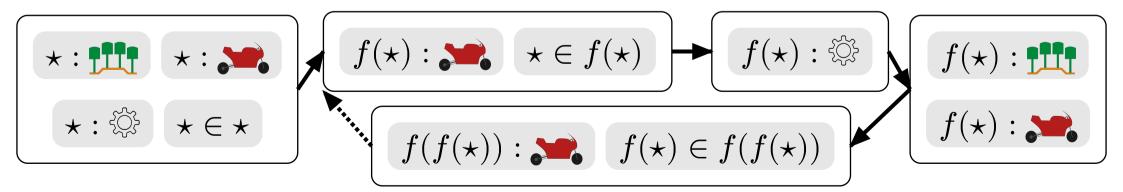
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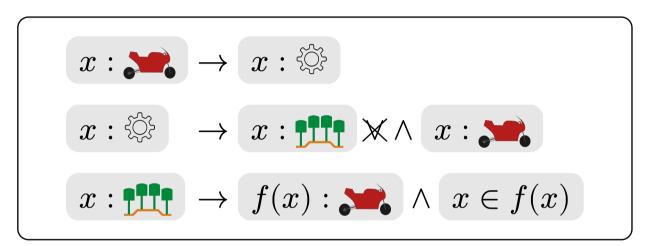


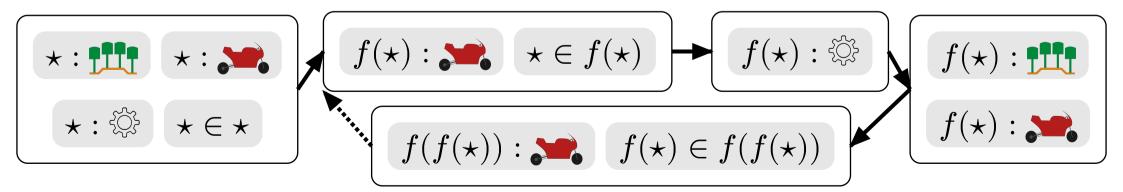
- 1. Replace Disjunctions
- 2. Use Critical Database
- 3. Run Chase
- 4. Stop on Cyclic Term





- 1. Replace Disjunctions
- 2. Use Critical Database
- 3. Run Chase
- 4. Stop on Cyclic Term
- 5. MFA iff no Cyclic Term

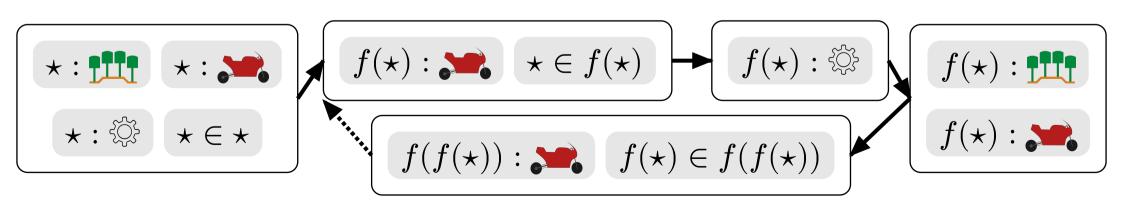


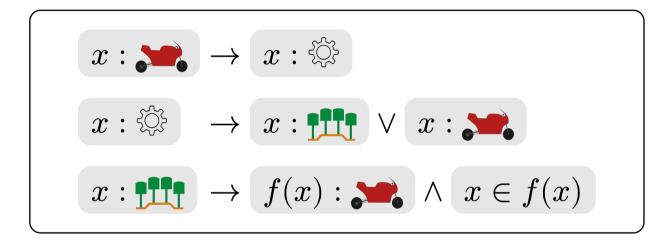


We want to know if the chase terminates on all databases for a rule set.

- 1. Replace Disjunctions
- 2. Use Critical Database
- 3. Run Chase
- 4. Stop on Cyclic Term
- 5. MFA iff no Cyclic Term

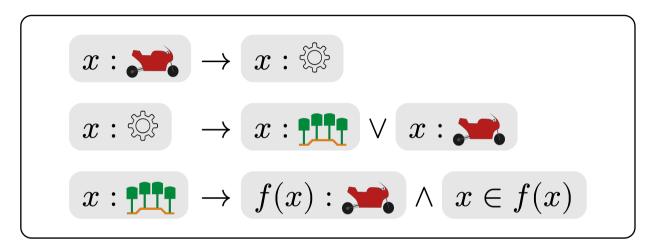
Theorem. If a rule set R is MFA, then R terminates on every database. Intuitively, the critical database chase subsumes every possible chase tree.





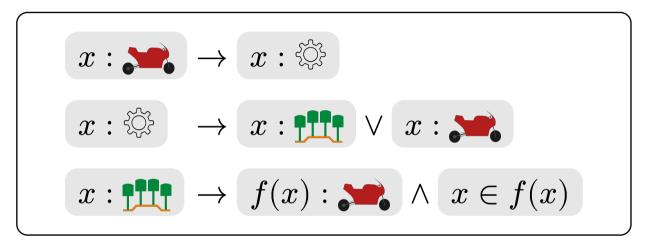
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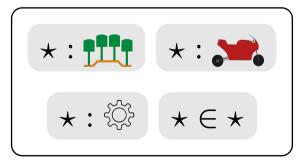
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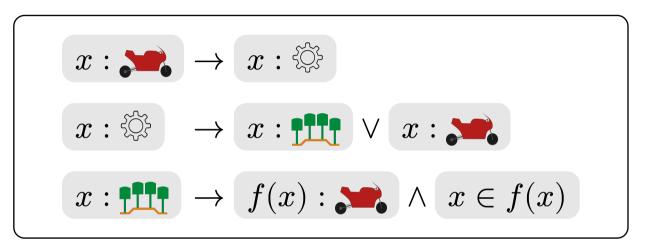
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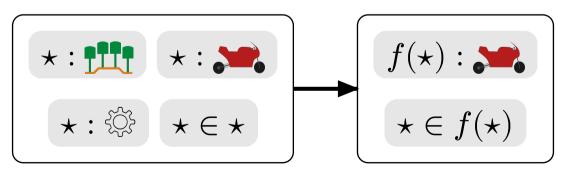
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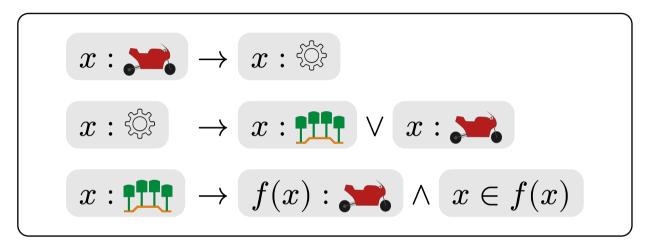


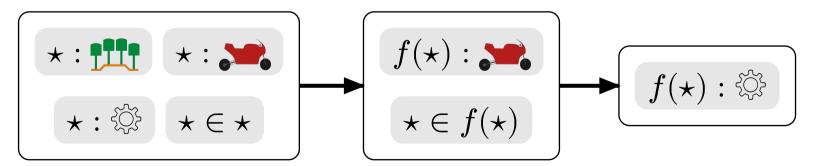
- 1. Use Critical Database
- Run Chase without blocked triggers



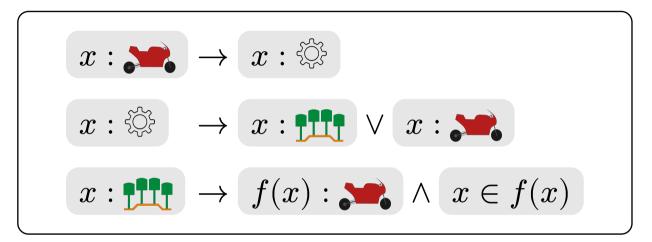


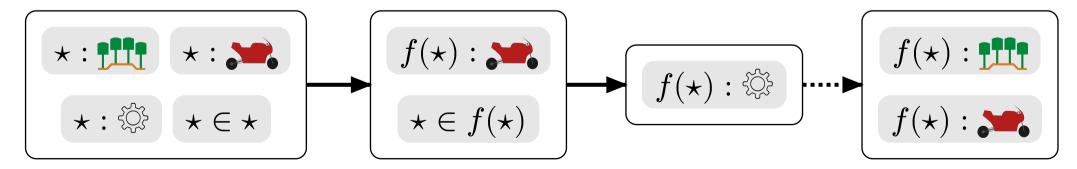
- 1. Use Critical Database
- 2. Run Chase without blocked triggers



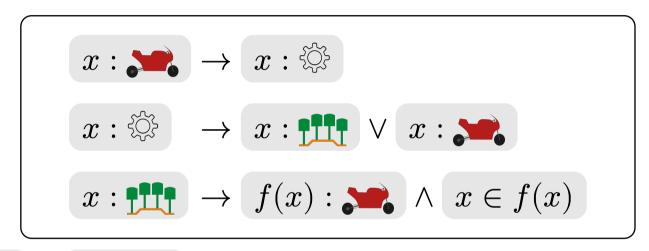


- 1. Use Critical Database
- Run Chase without blocked triggers





- 1. Use Critical Database
- Run Chase without blocked triggers



Is
$$x: \diamondsuit \to x:$$
 $\forall x:$ with $x/f(\star)$ blocked?

We want to know if the chase terminates on all databases for a rule set.

- 1. Use Critical Database
- Run Chase without blocked triggers

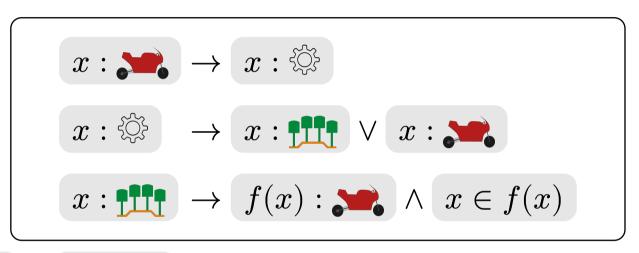
Is
$$x: \diamondsuit \to x: \longrightarrow with x/f(\star)$$
 blocked?

1. Backtrack facts for $f(\star)$

$$f(\star):$$

We want to know if the chase terminates on all databases for a rule set.

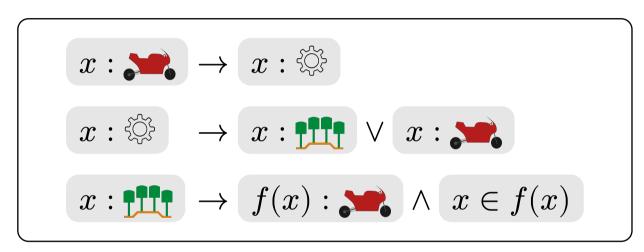
- 1. Use Critical Database
- Run Chase without blocked triggers



Is
$$x: \diamondsuit \to x: \circlearrowleft \lor x: \longleftrightarrow$$
 with $x/f(\star)$ blocked?

1. Backtrack facts for $f(\star)$

- 1. Use Critical Database
- Run Chase without blocked triggers



Is
$$x: \diamondsuit \to x: \square \lor x: \square$$
 with $x/f(\star)$ blocked?

- 1. Backtrack facts for $f(\star)$
- 2. Check Applicability



We want to know if the chase terminates on all databases for a rule set.

- 1. Use Critical Database
- 2. Run Chase without blocked triggers

Lemma. If a trigger is blocked it can never be applied in any chase tree.

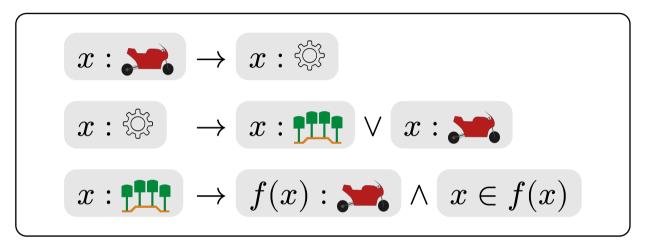
Intuitively, because the backtracked facts must occur in every chase tree.

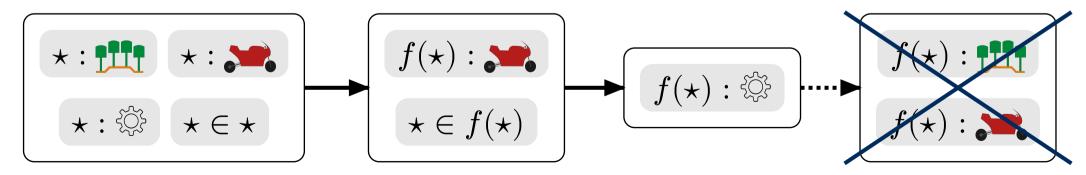
Is
$$x: \diamondsuit \to x: \mathbf{T} \lor x: \mathbf{W}$$
 with $x/f(\star)$ blocked? Yes!

- 1. Backtrack facts for $f(\star)$
- 2. Check Applicability

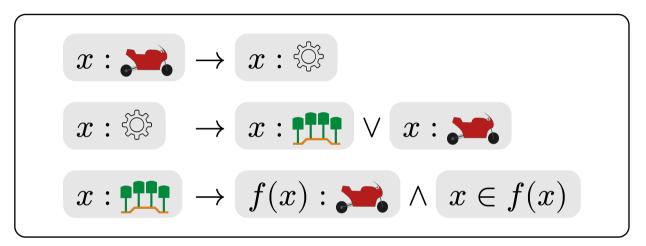


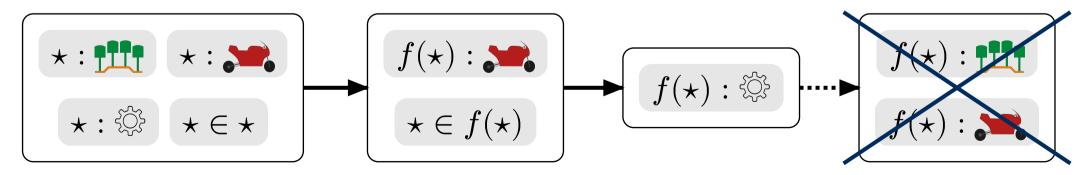
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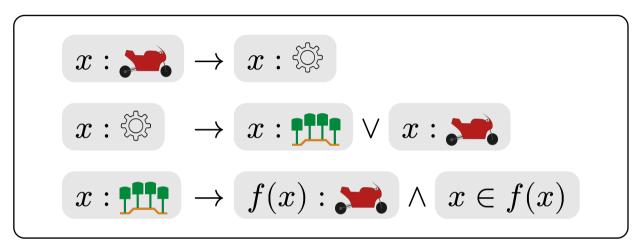


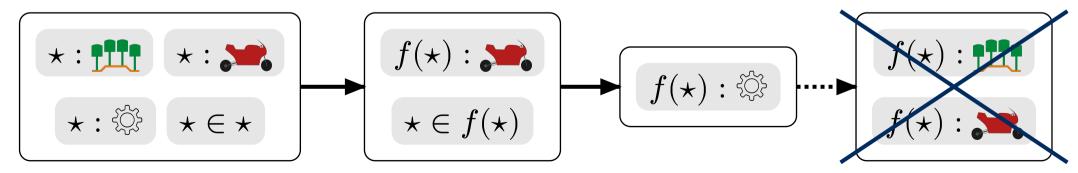
- 1. Use Critical Database
- Run Chase without blocked triggers
- 3. Stop on Cyclic Term





- 1. Use Critical Database
- Run Chase without blocked triggers
- 3. Stop on Cyclic Term
- 4. DMFA iff no Cyclic Term

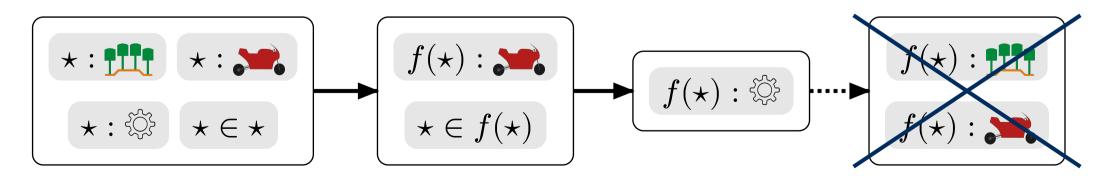


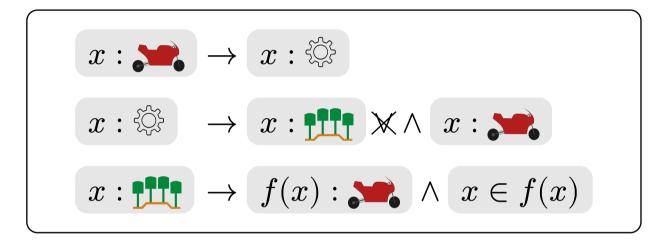


We want to know if the chase terminates on all databases for a rule set.

- 1. Use Critical Database
- Run Chasewithout blocked triggers
- 3. Stop on Cyclic Term
- 4. DMFA iff no Cyclic Term

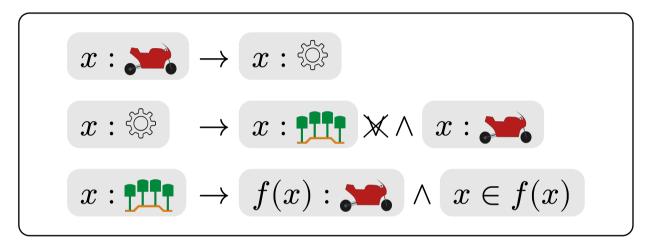
Theorem. If a rule set R is DMFA, then R terminates on every database. Again, the critical database chase subsumes every possible chase tree.



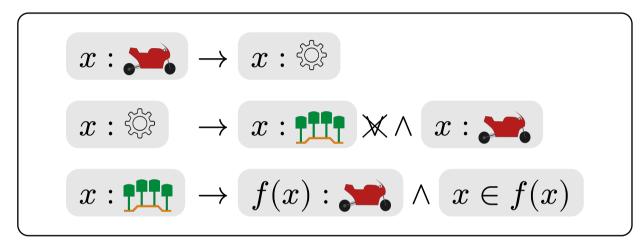


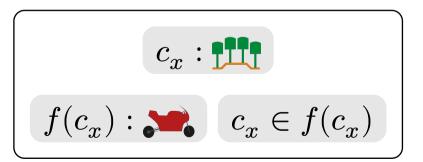
We want to know if the chase always loops on some DB for a rule set.

1. Ignore Disjunctive Rules

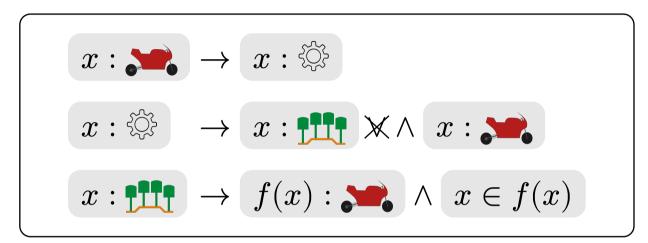


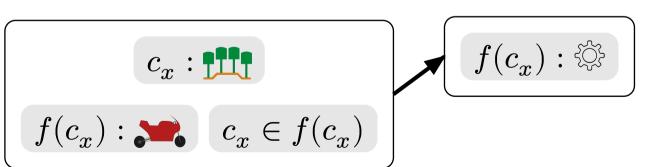
- 1. Ignore Disjunctive Rules
- 2. Use Rule Databases



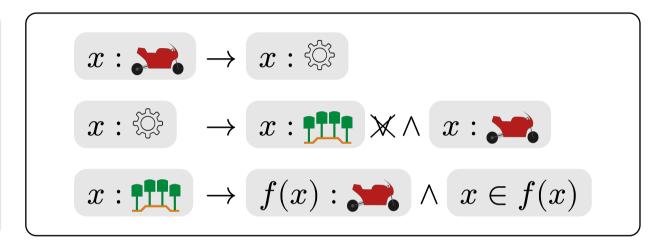


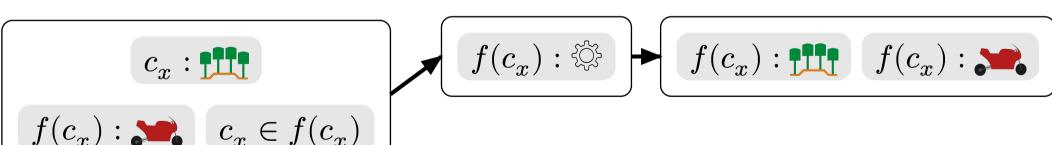
- 1. Ignore Disjunctive Rules
- 2. Use Rule Databases
- 3. Run Chase



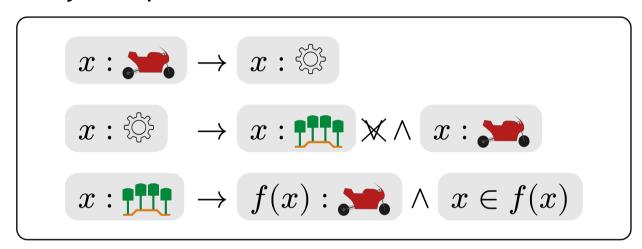


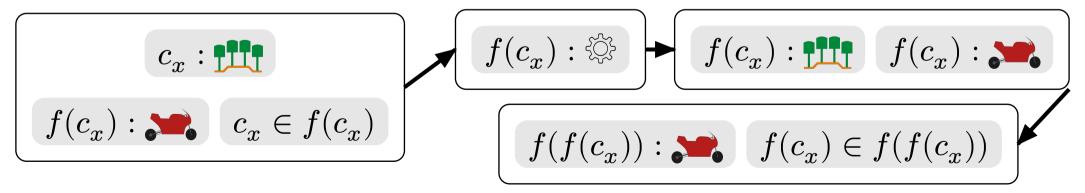
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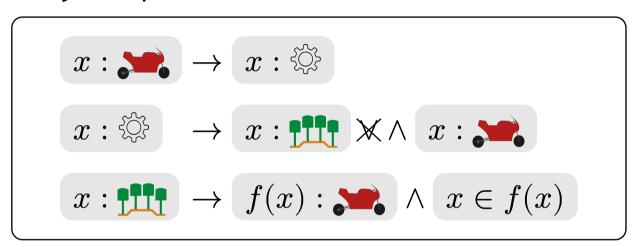


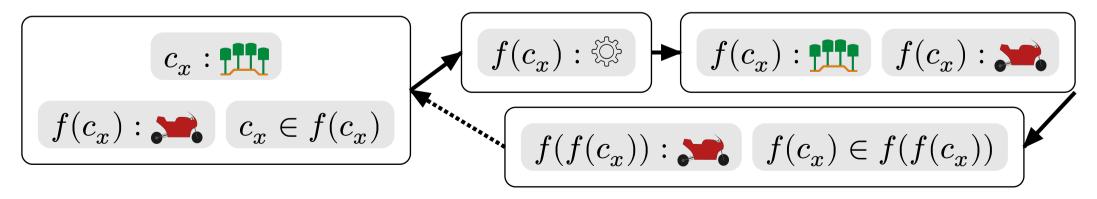
- 1. Ignore Disjunctive Rules
- 2. Use Rule Databases
- 3. Run Chase
- 4. Not beyond Cyclic Term





- 1. Ignore Disjunctive Rules
- 2. Use Rule Databases
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- 4. Not beyond Cyclic Term
- 5. MFC iff Rule-Cyclic Term

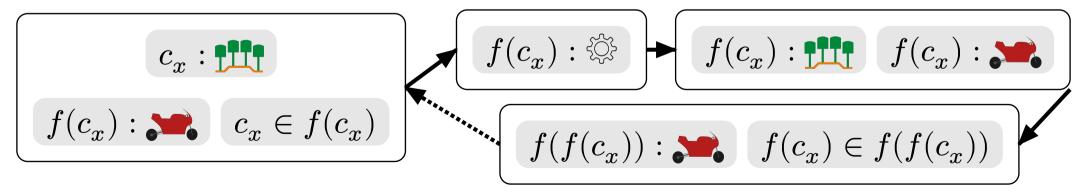




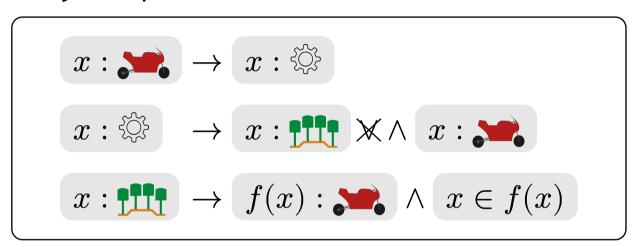
We want to know if the chase always loops on some DB for a rule set.

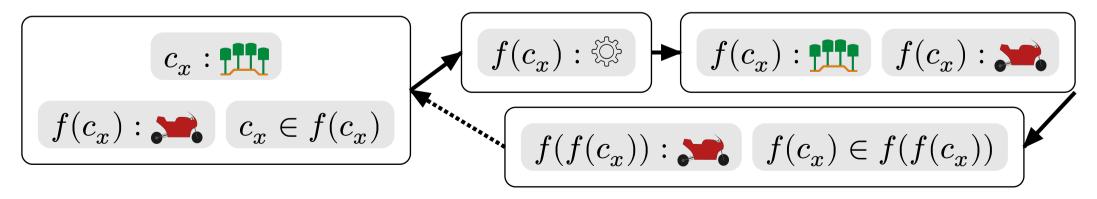
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Theorem. If a rule set R is MFC, then there is a rule database for a rule in R that only admits infinite chase trees. Intuitively, the loop occurs in every CT.

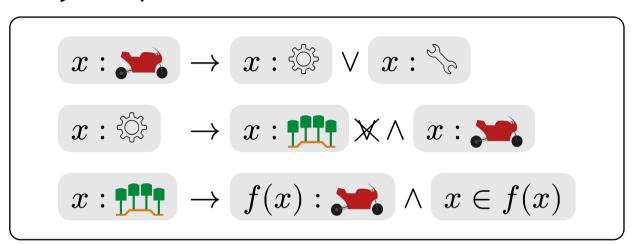


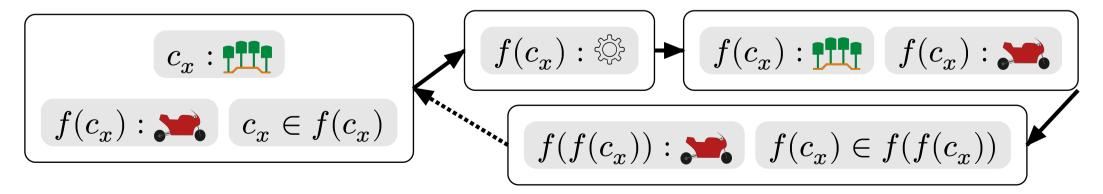
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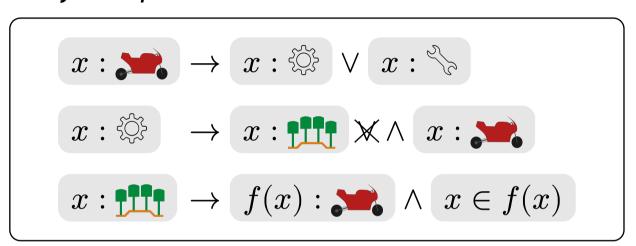


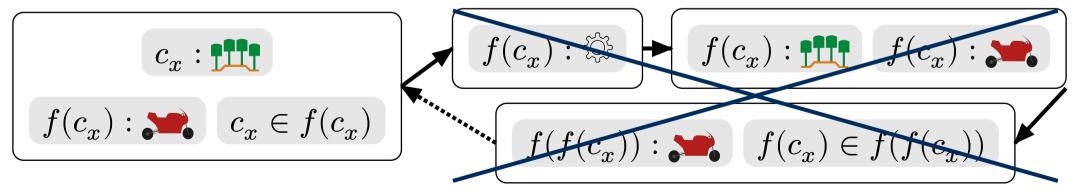


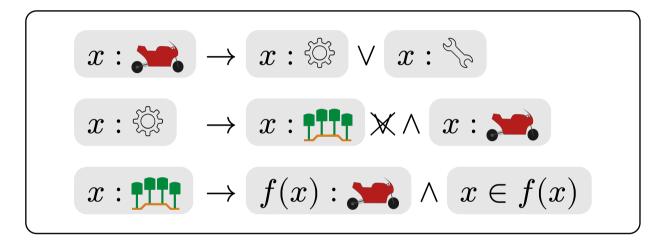
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- 4. Not beyond Cyclic Term
- 5. MFC iff Rule-Cyclic Term

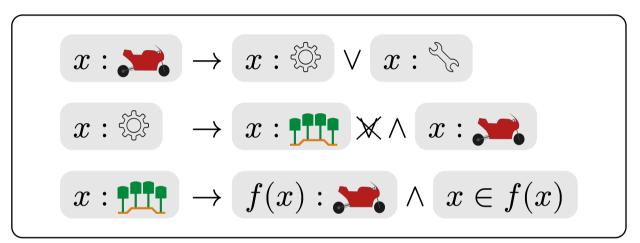






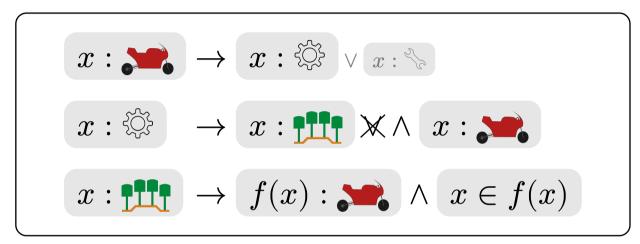
We want to know if the chase always loops on some DB for a rule set.

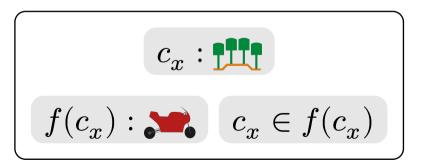
1. Ignore Disjunctive Rules



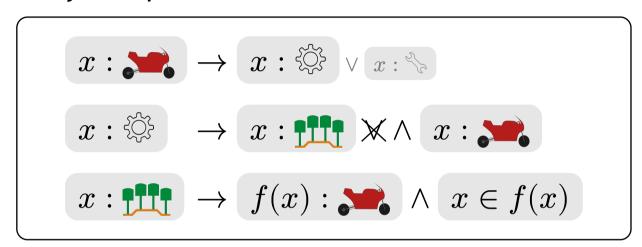
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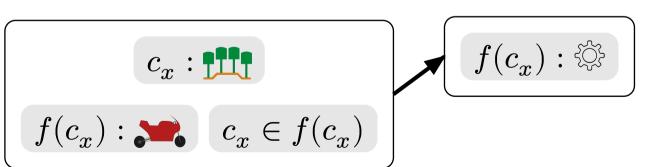
1. Rule DB and Head-Choice



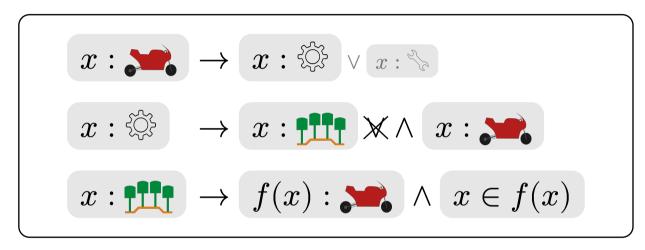


- 1. Rule DB and Head-Choice
- Run Chase only unblockable trgs





- 1. Rule DB and Head-Choice
- Run Chaseonly unblockable trgs



Is
$$x: \longrightarrow x: \diamondsuit \lor x: \diamondsuit$$
 with $x/f(c_x)$ unblockable?

We want to know if the chase always loops on some DB for a rule set.

- 1. Rule DB and Head-Choice
- 2. Run Chase only unblockable trgs

Lemma. If a trigger is unblockable, (1) its output must occur if its body is satisfied and (2) there are infinitely many *similar* unblockable triggers.

Is $x: \longrightarrow x: \diamondsuit \lor x: \diamondsuit$ with $x/f(c_x)$ unblockable?

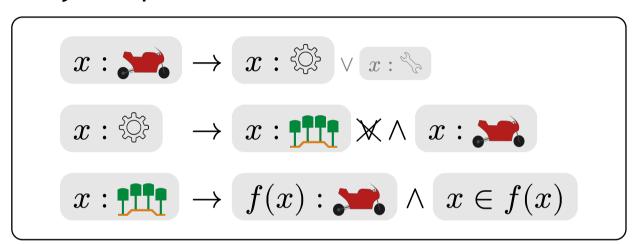
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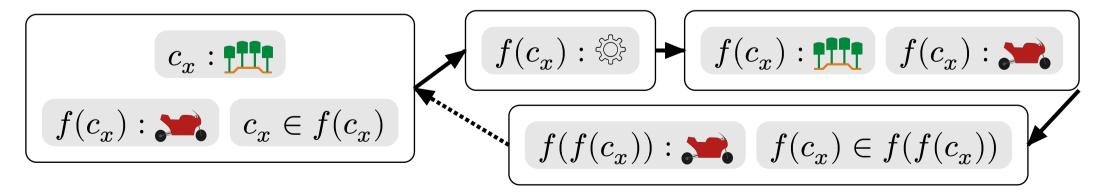
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Is $x: \longrightarrow x: \diamondsuit \lor x: \diamondsuit$ with $x/f(c_x)$ unblockable? Yes! I will explain on the next slide.

- 1. Rule DB and Head-Choice
- 2. Run Chase only unblockable trgs
- 3. Not beyond Cyclic Term
- 4. DMFC iff Rule-Cycl. Term

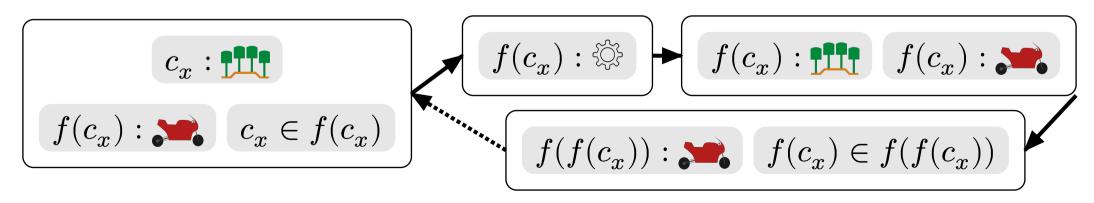




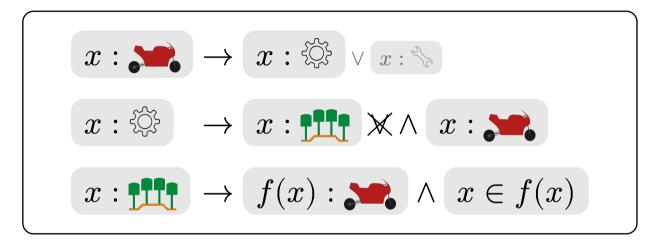
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- 1. Rule DB and Head-Choice
- Run Chaseonly unblockable trgs
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Theorem. If a rule set R is DMFC, then there is a rule database for a rule in R that only admits infinite chase trees. Infinite repetition of loop not obvious.









1. Backtrack facts for $f(c_x)$

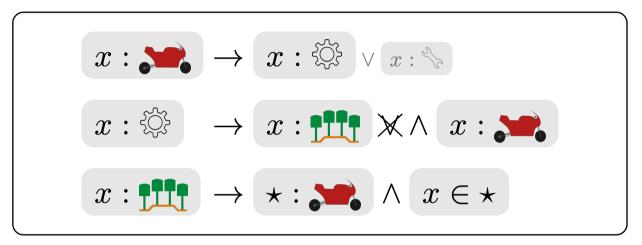
$$f(c_x): \ c_x \in f(c_x)$$



- 1. Backtrack facts for $f(c_x)$
- 2. Add critical facts

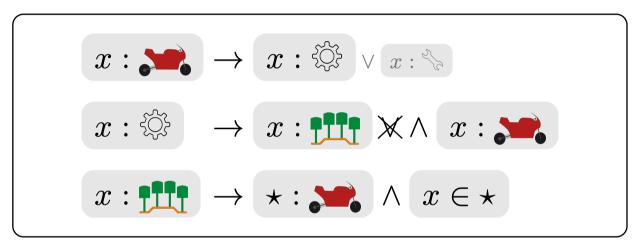


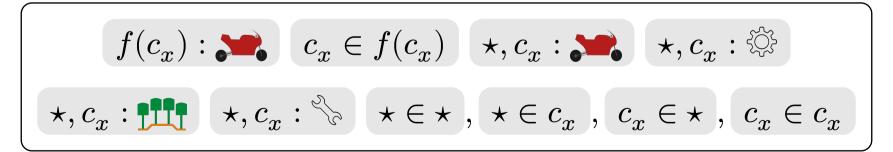
- 1. Backtrack facts for $f(c_x)$
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- 3. Chase with Star-Rules except checked trigger





- 1. Backtrack facts for $f(c_x)$
- 2. Add critical facts
- 3. Chase with Star-Rules except checked trigger
- 4. Unblk iff applicable



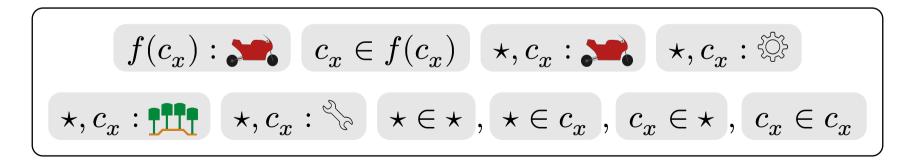




We compute an overapproximation of facts that can occur in the chase.

- 1. Backtrack facts for $f(c_x)$
- 2. Add critical facts
- 3. Chase with Star-Rules except checked trigger
- 4. Unblk iff applicable

Lemma. If a trigger is unblockable, (1) its output must occur if its body is satisfied and (2) there are infinitely many *similar* unblockable triggers.



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